

# A note on graph reconstruction

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## Abstract

Suppose  $G$  and  $G'$  are graphs on the same vertex set  $V$  such that for each  $x \in V$  there is an isomorphism  $\theta_x$  of  $G - x$  to  $G' - x$ . We prove in this paper that if there is a vertex  $x \in V$  and an automorphism  $\sigma$  of  $G - x$  such that  $\theta_x$  agrees with  $\sigma$  on all except for at most three vertices of  $V - x$ , then  $G$  is isomorphic to  $G'$ . As a corollary we prove that if a graph  $G$  has a vertex which is contained in at most three bad pairs, then  $G$  is reconstructible. Here a pair of vertices  $x, y$  of a graph  $G$  is called a bad pair if there exist  $u, v \in V(G)$  such that  $\{u, v\} \neq \{x, y\}$  and  $G - \{x, y\}$  is isomorphic to  $G - \{u, v\}$ .

All graphs discussed here are finite simple graphs. The vertex set and edge set of a graph  $G$  are denoted by  $V(G)$  and  $E(G)$  respectively. If  $A$  is a subset of  $V(G)$ , we use  $G|A$  and  $G - A$  to denote the subgraphs of  $G$  induced by  $A$  and  $V(G) - A$  respectively. When  $A = \{x\}$  is a singleton, we use  $G - x$  instead of  $G - \{x\}$ . For two subsets  $X, Y$  of  $V(G)$ , we use  $e_G(X, Y)$  to denote the number of edges joining a vertex of  $X$  to a vertex of  $Y$ . For brevity, we write  $e_G(x, X)$  for  $e_G(\{x\}, X)$  and  $e_G(X)$  for  $e_G(X, X)$ . The degree of  $x$  in  $G$  is denoted by  $d_G(x)$ .

We shall use some notations defined in [4]. Two graphs  $G$  and  $H$  are *hypomorphic* if there exists a bijection  $f : V(G) \rightarrow V(H)$  such that  $G - x$  is

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isomorphic to  $H - f(x)$  for each vertex  $x$  of  $G$ . Such a mapping  $f$  is called a *hypomorphism* of  $G$  to  $H$ . Obviously each isomorphism is a hypomorphism. The converse is not true. However the well-known reconstruction conjecture (cf. [3]) asserts that the existence of a hypomorphism of  $G$  to  $H$  implies the existence of an isomorphism of  $G$  to  $H$ . Let  $f : V(G) \rightarrow V(H)$  be a hypomorphism; and for each  $x \in V(G)$ , let  $p_x : G - x \rightarrow H - f(x)$  be an isomorphism. Define  $\theta_x = f^{-1}p_x$ , where mappings are composed from right to left. The mapping  $\theta_x$  is a permutation of  $V(G) - x$ , and does not act on  $x$ . We call  $\theta_x$  a *partial permutation* of  $V(G)$ .

Note that  $H - f(x) = f\theta_x(G - x)$ , so that  $H = \cup_x (H - f(x)) = f(\cup_x \theta_x(G - x))$ . The graph  $G' = \cup_x \theta_x(G - x)$  is called a *hypomorph* of  $G$ . A hypomorph  $G'$  of  $G$  is actually a graph on the same vertex set  $V = V(G)$  which is hypomorphic to  $G$  with identity being a hypomorphism. We note that  $G$  may have many hypomorphs, derived from different hypomorphisms. In particular  $G$  is a hypomorph of itself. If all hypomorphs of  $G$  are isomorphic then we say  $G$  is *reconstructible*.

W.L. Kocay [4] studied some basic properties of these partial permutations  $\theta_x$  as well as the partial automorphisms  $\theta_{xy} = \theta_x^{-1}\theta_y$ . Some sufficient conditions in terms of these mappings are given in [4] so that  $G$  is isomorphic to its hypomorph  $G'$ . It was shown in [4] that if there exist distinct vertices  $x, y \in V(G)$  such that  $\theta_x \in \text{Aut}(G - x)$  and  $\theta_y \in \text{Aut}(G - y)$  then  $G$  is isomorphic to  $G'$ .

In this note, we show that if there is a vertex  $x \in V(G)$  such that  $\theta_x$  is very "close" to an automorphism of  $G - x$ , then  $G$  is isomorphic to  $G'$ . To be precise, we will prove the following:

**Theorem 1** *Suppose  $G'$  is the hypomorph of  $G$  defined as above. If for a vertex  $x$  of  $G$ , there exists an automorphism  $\sigma \in \text{Aut}(G - x)$  of  $G - x$  which agrees with  $\theta_x$  on all except for at most three vertices of  $G - x$ , then  $G$  is isomorphic to  $G'$ .*

We call an unordered pair of vertices  $\{x, y\}$  of a graph  $G$  a *bad pair* if there exist  $u, v \in V(G)$  such that  $\{u, v\} \neq \{x, y\}$  and  $G - \{x, y\}$  is isomorphic to  $G - \{u, v\}$ . It was proved in [2] that if a graph  $G$  has a vertex  $x$  which

is contained in no bad pairs then  $G$  is reconstructible. As a consequence of Theorem 1, we obtain the following result:

**Corollary 1** *If  $G$  has a vertex which is contained in at most three bad pairs, then all hypomorphs of  $G$  are isomorphic to  $G$  and hence  $G$  is reconstructible.*

**Proof:** Suppose that  $G$  has a vertex  $x$  which is contained in at most three bad pairs. We need to show that any hypomorph  $G'$  of  $G$  is isomorphic to  $G$ .

For each  $x \in V$ , let  $\theta_x : V - x \rightarrow V - x$  be an isomorphism of  $G - x$  to  $G' - x$ . We now show that if  $\theta_x(y) \neq y$  for some vertex  $y \in V(G - x)$ , then  $\{x, y\}$  is a bad pair of  $G$ . Indeed, let  $z = \theta_x(y)$  and let  $w = \theta_z^{-1}(x)$ , where  $\theta_z : V - z \rightarrow V - z$  is an isomorphism of  $G - z$  to  $G' - z$ . Then  $G - \{x, y\}$  is isomorphic to  $G' - \{x, z\}$ , which is isomorphic to  $G - \{z, w\}$ . Since  $z \neq x$ , and  $z \neq y$ ,  $\{x, y\} \neq \{z, w\}$ , and hence  $\{x, y\}$  is a bad pair of  $G$ .

Because  $x$  is contained in at most three bad pair of  $G$ , we have that  $\theta_x(y) = y$  for at least  $|V(G)| - 4$  vertices  $y$  of  $G - x$ , i.e., the identity, which is an automorphism of  $G - x$ , agrees with  $\theta_x$  on all except for at most three vertices of  $G - x$ . Therefore  $G$  is isomorphic to  $G'$  by Theorem 1.

We now proceed to prove Theorem 1. The following lemma is an easy consequence of the fact that  $d_G(v) = d_{G'}(v)$  for all  $v \in V$  (cf. [1]).

**Lemma 1** *Suppose  $\sigma \in \text{Aut}(G - x)$  is an automorphism of  $G - x$ . If there is a subset  $A \subset V - \{x\}$  such that  $\sigma(A) = \theta_x(A)$ , then let  $\sigma(A) = B$ , we have  $e_G(x, B) = e_{G'}(x, B)$ .*

**Proof.** Since  $\sigma \in \text{Aut}(G - x)$  and  $\sigma(A) = B$ , we obtain  $e_G(B) = e_G(A) = e_{G'}(\theta_x(A)) = e_{G'}(B)$ . Similarly  $e_G(B, (V - x) - B) = e_{G'}(B, (V - x) - B)$ .

It is clear that  $e_G(B, V) = e_G(B) + e_G(B, (V - x) - B) + e_G(x, B)$ , and  $e_{G'}(B, V) = e_{G'}(B) + e_{G'}(B, (V - x) - B) + e_{G'}(x, B)$ . Also we have  $e_G(B, V) = \sum_{v \in B} d_G(v) - e_G(B) = \sum_{v \in B} d_{G'}(v) - e_{G'}(B) = e_{G'}(B, V)$ . Therefore  $e_G(x, B) = e_{G'}(x, B)$ . ■

**Corollary 2** *If  $v \in V(G - x)$  and  $\sigma(v) = u = \theta_x(v)$ , while  $\sigma \in \text{Aut}(G - x)$ , then  $(x, u) \in E(G)$  if and only if  $(x, u) \in E(G')$ .*

**Proof of Theorem 1:** Let  $x \in V$  be a vertex of  $G$  and let  $\sigma \in \text{Aut}(G - x)$  be an automorphism of  $G - x$  which agrees with  $\theta_x$  on all except for at most three vertices. We shall prove that  $G$  is isomorphic to  $G'$ .

If  $\sigma(v) = \theta_x(v)$  for all  $v \in V(G - x)$  then  $G$  and  $G'$  are identical. Indeed, for any edge  $(u, v) \in E(G)$  which does not contain  $x$  as an end point, we have  $(\sigma^{-1}(u), \sigma^{-1}(v)) \in E(G)$ . Therefore  $(\theta_x(\sigma^{-1}(u)), \theta_x(\sigma^{-1}(v))) \in E(G')$ . But  $\theta_x = \sigma$ , so  $(u, v) \in E(G')$ . For an edge  $(x, u) \in E(G)$  which does contain  $x$  as an end point, we have  $(x, u) \in E(G')$  by Corollary 2. Thus  $G$  is isomorphic to  $G'$ .

Next we consider the case that there are exactly two vertices, say  $v_1, v_2$ , of  $G - x$ , such that  $\sigma(v_i) \neq \theta_x(v_i)$  ( $i = 1, 2$ ). Let  $u_i = \sigma(v_i)$  for  $i = 1, 2$ . Then we must have  $\theta_x(v_1) = u_2$  and  $\theta_x(v_2) = u_1$ . By Lemma 1, we have  $(x, v) \in E(G)$  if and only if  $(x, v) \in E(G')$  for all vertices  $v \in V - x$  not equal to  $u_1$  or  $u_2$ ; and  $e_G(x, \{u_1, u_2\}) = e_{G'}(x, \{u_1, u_2\})$ . If  $e_G(x, \{u_1, u_2\}) = 2$  or  $0$ , then it is easy to show (similar to the argument in the previous paragraph) that the mapping  $g : V \rightarrow V$  defined by  $g = \theta_x \sigma^{-1}$  on  $V - x$  and  $g(x) = x$  is an isomorphism of  $G$  to  $G'$ .

Thus we assume that  $e(x, \{u_1, u_2\}) = 1$ . Without loss of generality we assume that  $(x, u_1) \in E(G)$ . If  $(x, u_2) \in E(G')$  then again the mapping  $g$  defined in the previous paragraph is an isomorphism of  $G$  to  $G'$ . Thus we assume that  $(x, u_1) \in E(G')$ . We claim that in this case  $G$  is identical to  $G'$ , i.e., the identity is an isomorphism of  $G$  to  $G'$ .

Otherwise there are vertices  $a, b \in V$  such that  $(a, b) \in E(G)$  and  $(a, b) \notin E(G')$ . It is easy to verify (similar to the proof of the case  $\sigma = \theta_x$ ) that  $G$  and  $G'$  are identical on  $V - \{u_1, u_2\}$ , and they are also identical on  $\{x, u_1, u_2\}$ . Thus we may assume that  $a \in \{u_1, u_2\}$  and  $b \in V - \{x, u_1, u_2\}$ .

Without loss of generality, we can assume that  $a = u_1$ . (If  $a = u_2$  then we have the equivalent of  $a = u_1$  in  $G'$  and we may interchange the roles of  $G$  and  $G'$ .) Since  $G - b$  is isomorphic to  $G' - b$ , these two graphs have the same degree sequence. However for all vertices  $v \neq u_1, u_2$ , we have  $d_{G-b}(v) = d_{G'-b}(v)$  (as  $d_G(v) = d_{G'}(v)$  and  $(b, v) \in E(G)$  if and only if  $(b, v) \in E(G')$ ). Therefore we must have  $\{d_{G-b}(u_1), d_{G-b}(u_2)\} = \{d_{G'-b}(u_1), d_{G'-b}(u_2)\}$ .

Since  $(u_1, b) \in E(G)$ , we have  $(\sigma^{-1}(u_1), \sigma^{-1}(b)) \in E(G)$  and hence  $(\theta_x \sigma^{-1}(u_1), \theta_x \sigma^{-1}(b)) \in E(G')$ , i.e.,  $(u_2, b) \in E(G')$ . Similarly,  $(u_1, b) \notin$

$E(G')$  implies  $(u_2, b) \notin E(G)$ . Therefore  $d_G(u_1) = d_{G-b}(u_1) + 1$ ,  $d_G(u_2) = d_{G-b}(u_2)$ ,  $d_{G'}(u_1) = d_{G'-b}(u_1)$ , and  $d_{G'}(u_2) = d_{G'-b}(u_2) + 1$ .

As  $d_G(u_1) = d_{G'}(u_1)$ ,  $d_G(u_2) = d_{G'}(u_2)$ , we conclude that  $d_G(u_1) = d_G(u_2) = d_{G'}(u_1) = d_{G'}(u_2)$ , which implies that  $d_{G-x}(u_1) \neq d_{G'-x}(u_2)$ . This is a contradiction, as  $\theta_x \sigma^{-1}$  is an isomorphism of  $G-x$  to  $G'-x$  which sends  $u_1$  to  $u_2$ . Therefore  $G$  and  $G'$  are identical.

Finally we consider the case that there are three vertices of  $G-x$ , say  $v_1, v_2, v_3$ , such that  $\sigma(v_i) \neq \theta_x(v_i)$  ( $i = 1, 2, 3$ ). Without loss of generality, we may assume that  $\theta_x(v_1) = \sigma(v_2) = u_2$ ,  $\theta_x(v_2) = \sigma(v_3) = u_3$  and  $\theta_x(v_3) = \sigma(v_1) = u_1$ .

By Lemma 1,  $(x, v) \in E(G)$  if and only if  $(x, v) \in E(G')$  for all vertices  $v$  of  $G-x$  not equal to  $u_1, u_2$  or  $u_3$ ; and  $e_G(x, \{u_1, u_2, u_3\}) = e_{G'}(x, \{u_1, u_2, u_3\})$ . If  $e_G(x, \{u_1, u_2, u_3\}) = 3$  or  $0$ , then again it is easy to verify that the mapping  $g : V \rightarrow V$  defined as  $g = \theta_x \sigma^{-1}$  on  $V-x$  and  $g(x) = x$  is an isomorphism of  $G$  to  $G'$ .

We now consider the case that  $e_G(x, \{u_1, u_2, u_3\}) = 1$ . The case  $e_G(x, \{u_1, u_2, u_3\}) = 2$  will follow easily by considering the complement graphs.

Without loss of generality, we assume that  $(x, u_1) \in E(G)$ . If  $(x, u_2) \in E(G')$  then again the mapping  $g$  defined above is an isomorphism of  $G$  to  $G'$ . We now assume that  $(x, u_2) \notin E(G')$ . Thus we have either  $(x, u_1) \in E(G')$  or  $(x, u_3) \in E(G')$ .

**Case 1:** Suppose  $(x, u_1) \in E(G')$ . We shall show that in this case the mapping  $g$  which sends  $u_2$  to  $u_3$ , sends  $u_3$  to  $u_2$ , and fixes every other vertices of  $V$  is an isomorphism of  $G$  to  $G'$ . Suppose  $d_G(u_1) = k + 1$ . then it is easy to see that  $d_{G-x}(u_2) = d_G(u_2) = d_{G-x}(u_3) = d_G(u_3) = k$ . Let  $S = V - \{u_1, u_2, u_3\}$ . For any vertex  $u \in S$ , we have  $\{d_{G-u}(u_1), d_{G-u}(u_2), d_{G-u}(u_3)\} = \{d_{G'-u}(u_1), d_{G'-u}(u_2), d_{G'-u}(u_3)\}$ , because  $G-u$  is isomorphic to  $G'-u$  (hence these two graphs have the same degree sequence), and  $d_{G-u}(v) = d_{G'-u}(v)$  for all vertices  $v \neq u_1, u_2, u_3$  (as  $d_G(v) = d_{G'}(v)$  and  $(u, v) \in E(G)$  if and only if  $(v, u) \in E(G')$ ). This implies that for any vertex  $u \in S$ , we have  $(u_1, u) \in E(G)$  if and only if  $(u_1, u) \in E(G')$ . Indeed if  $(u_1, u) \notin E(G)$  and  $(u_1, u) \in E(G')$ , then  $d_{G-u}(u_1) = k + 1$  and  $d_{G'-u}(u_i) \leq k$  for all  $i = 1, 2, 3$ ; and hence  $\{d_{G-u}(u_1), d_{G-u}(u_2), d_{G-u}(u_3)\} \neq \{d_{G'-u}(u_1), d_{G'-u}(u_2), d_{G'-u}(u_3)\}$ . Similar

contradiction can be derived if  $(u_1, u) \in E(G)$  and  $(u_1, u) \notin E(G')$ .

Recall that  $\theta_x \sigma^{-1}(u_3) = u_1$  and  $\theta_x \sigma^{-1}(u_1) = u_2$ , we conclude that for any  $u \in S - \{x\}$ ,  $(u, u_3) \in E(G)$  if and only if  $(u, u_1) \in E(G')$  if and only if  $(u, u_1) \in E(G)$  if and only if  $(u, u_2) \in E(G')$ . Furthermore  $(x, u_3) \notin E(G)$  and  $(x, u_2) \notin E(G')$ . Thus for all  $u \in S$ , we have  $(u, u_3) \in E(G)$  if and only if  $(u, u_2) \in E(G')$ . To prove that the mapping  $g$  defined above is an isomorphism of  $G$  to  $G'$ , it remains to show that  $(u_i, u_j) \in E(G)$  if and only if  $(g(u_i), g(u_j)) \in E(G')$  for  $i, j \in \{1, 2, 3\}$ .

First observe that because  $G - u_1$  is isomorphic to  $G' - u_1$ , and  $d_{G-u_1}(u) = d_{G'-u_1}(u)$  for any  $u \in S$ , we have  $\{d_{G-u_1}(u_2), d_{G-u_1}(u_3)\} = \{d_{G'-u_1}(u_2), d_{G'-u_1}(u_3)\}$ .

We now consider two subcases:

**Case 1(a):** Suppose that  $(u_1, u_2) \in E(G)$ . Then  $(u_2, u_3) \in E(G')$ , as  $\theta_x \sigma^{-1}(u_1) = u_2$  and  $\theta_x \sigma^{-1}(u_2) = u_3$ . Since  $d_{G-u_1}(u_2) = k - 1 \in \{d_{G'-u_1}(u_2), d_{G'-u_1}(u_3)\}$ , we must have  $(u_1, u_2) \in E(G')$  or  $(u_1, u_3) \in E(G')$ .

If  $(u_1, u_2) \in E(G')$ , then  $(u_1, u_3) \in E(G)$ , and hence  $\{d_{G-u_1}(u_2), d_{G-u_1}(u_3)\} = \{k - 1, k - 1\}$ . In order that  $\{d_{G'-u_1}(u_2), d_{G'-u_1}(u_3)\} = \{k - 1, k - 1\}$ , we must have  $(u_1, u_3) \in E(G')$ . This implies that  $(u_2, u_3) \in E(G)$ , and hence  $\{u_1, u_2, u_3\}$  induces a complete graph in both graphs  $G$  and  $G'$ .

If  $(u_1, u_2) \notin E(G')$  then  $(u_1, u_3) \in E(G')$ . This implies that  $(u_2, u_3) \in E(G)$  and  $(u_1, u_3) \notin E(G)$ , as  $\theta_x \sigma^{-1}$  is an isomorphism of  $G$  to  $G'$ .

In any case the restriction of  $g$  to  $\{u_1, u_2, u_3\}$  is an isomorphism.

**Case 1(b):** Suppose that  $(u_1, u_2) \notin E(G)$ . Then  $(u_2, u_3) \notin E(G')$ . If  $(u_1, u_3) \in E(G)$  then  $(u_1, u_2) \in E(G')$ . This implies that  $(u_1, u_3) \notin E(G')$  for otherwise we would have  $\{d_{G-u_1}(u_2), d_{G-u_1}(u_3)\} = \{k, k - 1\}$  and  $\{d_{G'-u_1}(u_2), d_{G'-u_1}(u_3)\} = \{k - 1, k - 1\}$ . Since  $\theta_x \sigma^{-1}(u_2) = u_3$  and  $\theta_x \sigma^{-1}(u_3) = u_1$ , we know that  $(u_2, u_3) \notin E(G)$ . Thus the restriction of  $g$  to  $\{u_1, u_2, u_3\}$  is an isomorphism.

If  $(u_1, u_3) \notin E(G)$  then  $(u_1, u_2) \notin E(G')$ . This implies that  $(u_1, u_3) \notin E(G')$ , for otherwise we would have  $\{d_{G-u_1}(u_2), d_{G-u_1}(u_3)\} = \{k, k\}$  and  $\{d_{G'-u_1}(u_2), d_{G'-u_1}(u_3)\} = \{k, k - 1\}$ . Thus again we have  $(u_2, u_3) \notin E(G)$ ,

and therefore the restriction of  $g$  to  $\{u_1, u_2, u_3\}$  is an isomorphism.

**Case 2:** Suppose that  $(x, u_3) \in E(G')$ . We shall show that in this case the mapping  $g$  which sends  $u_1$  to  $u_3$ , sends  $u_3$  to  $u_1$ , and fixes every other vertex of  $V$ , is an isomorphism of  $G$  to  $G'$ . The proof is very similar to that of Case 1 and we omit some details.

Let  $k = d_{G-x}(u_1)$ . Then  $d_G(u_1) = d_{G'}(u_1) = d_G(u_3) = d_{G'}(u_3) = k + 1$  and  $d_G(u_2) = d_{G'}(u_2) = k$ .

Similar to the argument in the proof of Case 1, we can show that for any vertex  $u \in S = V - \{u_1, u_2, u_3\}$ , we have  $\{d_{G-u}(u_1), d_{G-u}(u_2), d_{G-u}(u_3)\} = \{d_{G'-u}(u_1), d_{G'-u}(u_2), d_{G'-u}(u_3)\}$ . This implies that for any vertex  $u \in S$ , we have  $(u_2, u) \in E(G)$  if and only if  $(u_2, u) \in E(G')$  (cf. the proof of Case 1).

It remains to show that  $g$  restricted to  $\{u_1, u_2, u_3\}$  is an isomorphism. Similarly  $\{d_{G-u_2}(u_1), d_{G-u_2}(u_3)\} = \{d_{G'-u_2}(u_1), d_{G'-u_2}(u_3)\}$ , because  $G - u_2$  is isomorphic to  $G' - u_2$ , and  $d_{G-u_2}(u) = d_{G'-u_2}(u)$  for any  $u \in S$ . Again we consider two subcases:

**Case 2(a):** Suppose that  $(u_2, u_3) \in E(G)$ . Then  $(u_1, u_3) \in E(G')$ . This implies that  $(u_1, u_2) \in E(G')$  for otherwise  $\{d_{G-u_2}(u_1), d_{G-u_2}(u_3)\} \neq \{d_{G'-u_2}(u_1), d_{G'-u_2}(u_3)\}$ . This then implies that  $(u_1, u_3) \in E(G)$ , and hence the restriction of  $g$  to  $\{u_1, u_2, u_3\}$  is an isomorphism.

**Case 2(b):** Suppose that  $(u_2, u_3) \notin E(G)$ . Then  $(u_1, u_3) \notin E(G')$ . Similarly by using the condition that  $\{d_{G-u_2}(u_1), d_{G-u_2}(u_3)\} = \{d_{G'-u_2}(u_1), d_{G'-u_2}(u_3)\}$ , we can show that  $g$  restricted to  $\{u_1, u_2, u_3\}$  is an isomorphism. This completes the proof of Theorem 1.

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