

# Circular chromatic index of Cartesian products of graphs

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## Abstract

The *circular chromatic index* of a graph  $G$ , written  $\chi'_c(G)$ , is the minimum  $r$  permitting a function  $f: E(G) \rightarrow [0, r)$  such that  $1 \leq |f(e) - f(e')| \leq r - 1$  whenever  $e$  and  $e'$  are incident. Let  $G = H \square C_{2m+1}$ , where  $\square$  denotes Cartesian product and  $H$  is an  $(s - 2)$ -regular graph of odd order, with  $s \equiv 0 \pmod{4}$  (thus  $G$  is  $s$ -regular). We prove that  $\chi'_c(G) \geq s + \lfloor \lambda(1 - 1/s) \rfloor^{-1}$ , where  $\lambda$  is the minimum, over all bases of the cycle space of  $H$ , of the maximum length of a cycle in the basis. When  $H = C_{2k+1}$  and  $m$  is large, the lower bound is sharp. In particular, if  $m \geq 3k + 1$ , then  $\chi'_c(C_{2k+1} \square C_{2m+1}) = 4 + \lceil 3k/2 \rceil^{-1}$ , independent of  $m$ .

## 1 Introduction

The *chromatic index*  $\chi'(G)$  of a graph  $G$  is the minimum number of colors needed to color the edges so that incident edges receive distinct colors. In the case of a simple graph  $G$  (no loops or multiple edges), the famous theorem of Vizing [10] and Gupta [4] yields  $\Delta(G) \leq \chi'(G) \leq \Delta(G) + 1$ , where  $\Delta(G)$  is the maximum vertex degree in  $G$ .

With only two values available, it is common to say that a graph  $G$  is *Class 1* if  $\chi'(G) = \Delta(G)$  and *Class 2* otherwise. In this paper we consider a refinement of the chromatic index called the “circular chromatic index”. It equals  $\chi'(G)$  when  $G$  is Class 1, and otherwise it lies between  $\Delta(G)$  and  $\chi'(G)$ . To define it, we first describe a vertex coloring parameter.

Given a graph  $G$  and a real number  $r$ , an  $r$ -*coloring* of  $G$  is a function  $f: V(G) \rightarrow [0, r)$  such that  $1 \leq |f(x) - f(y)| \leq r - 1$  whenever  $x$  and  $y$  are adjacent. In essence, the set of colors form a circle of circumference  $r$ , and the colors assigned to adjacent vertices must differ by at least 1 (in each direction) along the circle.

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The *circular chromatic number* of  $G$ , written  $\chi_c(G)$ , is the infimum of all  $r$  such that  $G$  admits an  $r$ -coloring (the infimum can be replaced with minimum). There are many equivalent formulations of  $\chi_c(G)$  (see [12, 13] for surveys and many basic results). The definition here is not the most common but is useful for our results. Due to the elementary result that  $\chi(G) - 1 < \chi_c(G) \leq \chi(G)$  [9], the parameter  $\chi_c$  is a refinement of  $\chi$ , and this has motivated its extensive study over the past decade.

For a graph  $G$ , the *line graph*  $L(G)$  is the graph with vertex set  $E(G)$  whose adjacency relation is the incidence relation for edges in  $G$ . The *circular chromatic index*  $\chi'_c(G)$  is defined by  $\chi'_c(G) = \chi_c(L(G))$ . That is, we seek the smallest  $r$  permitting an  $r$ -coloring of the edges of  $G$ . Since  $\chi'(G) = \chi(L(G))$ , we have  $\chi'(G) - 1 < \chi'_c(G) \leq \chi'(G)$ , and  $\chi'_c$  is a refinement of  $\chi'$ . From the definition,  $\chi'_c(G)$  is at least the maximum number of pairwise incident edges. Thus  $\chi'_c(G) = \chi'(G)$  when  $G$  is Class 1. Otherwise,  $\Delta(G) < \chi'_c(G) \leq \Delta(G) + 1$ .

Several papers have been published about  $\chi'_c$ . It was proved in [2] that all 2-edge-connected graphs with maximum degree at most 3 have circular chromatic index at most  $11/3$ , except for two small graphs with circular chromatic index 4. In [5], it was proved that 2-edge-connected 3-regular graphs of large girth have circular chromatic index close to 3. This result was generalized in [6]: for any positive integer  $d$ , graphs with maximum degree  $d$  have circular chromatic index arbitrarily close to  $d$  if their girth is sufficiently large.

In this paper, we study the behavior of circular chromatic index under a product operation. Given graphs  $G$  and  $H$ , the *Cartesian product*  $G \square H$  is the graph with vertex set  $V(G) \times V(H)$  defined by making the pair  $(u, v)$  adjacent to the pair  $(u', v')$  if (1)  $u = u'$  and  $vv' \in E(H)$ , or (2)  $v = v'$  and  $uu' \in E(G)$ . It has long been known that  $\chi(G \square H) = \max\{\chi(G), \chi(H)\}$  [1, 8, 11]. The argument holds as well for  $\chi_c$ , so the behavior of  $\chi_c$  is trivial under the Cartesian product.

The behavior of  $\chi'_c$  is more interesting. If  $G \square H$  is Class 1, then  $\chi'_c(G \square H) = \Delta(G \square H)$ , so we consider only products that are Class 2. The product is Class 1 when  $G$  or  $H$  is Class 1 [7] or when  $G$  and  $H$  both have perfect matchings [7]. To avoid Class 1, let  $G$  and  $H$  be regular graphs with odd order. The product  $G \square H$  is then also regular with odd order, and a regular graph is Class 1 if and only if it has an edge-coloring in which every color class is a perfect matching, which does not exist in  $G \square H$ .

In particular, we consider the product of an odd cycle with a regular graph  $H$  of odd order, where the degree of the vertices in  $H$  is congruent to 2 modulo 4. We prove that  $\chi'_c(H \square C_{2m+1}) \geq s + \lfloor \lambda(1 - 1/s) \rfloor^{-1}$ , where  $\lambda$  is the maximum length of the cycles in some basis of the cycle space of  $H$  (choosing the basis to make  $\lambda$  smallest gives the best lower bound). We also prove that the bound is sharp when  $H$  is an odd cycle and  $m$  is large. Indeed,  $\chi'_c(H \square C_{2m+1})$  always decreases to a limit as  $m$  increases. In particular, if  $m \geq 3k + 1$ , then  $\chi'_c(C_{2k+1} \square C_{2m+1}) = 4 + \lceil 3k/2 \rceil^{-1}$ , independent of  $m$ .

## 2 Properties of $r$ -Colorings

We view the color set  $[0, r)$  for a  $r$ -coloring of a graph as the set of real numbers modulo  $r$ . Thus we interpret it as a circle  $C^r$  of circumference  $r$ , by identifying 0 and  $r$ . For  $a, b \in C^r$ , we write  $[a, b]_r$  for the set in  $C^r$  moving from  $a$  to  $b$  through increasing values. That is,  $[a, b]_r = [a, b]$  when  $a \leq b$ , while  $[a, b]_r = [a, r) \cup [0, b]$  when  $a > b$ . For convenience, we extend this notation to all real numbers  $a$  and  $b$  by letting  $[a, b]_r = [a \bmod r, b \bmod r]_r$ , where  $a \bmod r$  and  $b \bmod r$  are the remainders of  $a$  and  $b$  upon division by  $r$ . The intervals  $[a, b)_r$ ,  $(a, b]_r$  and  $(a, b)_r$  are defined similarly. We use  $\ell([a, b])$  to denote the length of the interval  $[a, b]$ , and we define a measure of distance on the circle as  $|a - b|_r = \min\{\ell([a, b]_r), \ell([b, a]_r)\}$ . An  $s$ -clique is a set of  $s$  pairwise adjacent vertices.

**Lemma 2.1** *Let  $G$  be a graph and  $f$  be an  $r$ -coloring of  $G$ , where  $r = s + \epsilon$  with  $s \in \mathbb{N}$  and  $\epsilon < 1/2$ . If  $Q$  is an  $s$ -clique in  $G$  and  $v \in Q$ , then each set  $[f(v) + i, f(v) + i + \epsilon]_r$  for  $0 \leq i \leq s - 1$  contains the color of exactly one vertex in  $Q$ . If  $X$  and  $Y$  are intersecting  $s$ -cliques, then for each  $x \in X$  there is a unique  $y \in Y$  such that  $|f(y) - f(x)|_r \leq \epsilon$ .*

**Proof.** Since the colors on vertices of  $Q$  must pairwise differ by at least 1, the  $i$ th such color after  $f(v)$  must be at least  $i$  units later along the circle. It cannot be more than  $i + \epsilon$  units later, since  $s - i$  subsequent colors are encountered in returning to  $f(v)$ .

Now consider  $v \in X \cap Y$ . With  $x_0 = y_0 = v$ , let  $x_i$  be the  $i$ th vertex of  $X$  whose color is encountered moving upward from  $f(v)$  around the circle (similarly define  $y_i$ ). By the preceding paragraph, both  $f(x_i)$  and  $f(y_i)$  lie in  $[f(v) + i, f(v) + i + \epsilon]_r$ , for  $1 \leq i \leq s - 1$ . Hence they differ by at most  $\epsilon$ . Furthermore, since  $\epsilon < 1/2$ , the distance between two such intervals is more than  $\epsilon$ , so  $y_i$  is the only vertex of  $Y$  whose color is within  $\epsilon$  of  $f(x_i)$ . ■

To facilitate proofs, we interpret vertex colorings as edge-weightings of orientations. Let  $\vec{G}$  be an orientation of a graph  $G$ . For a weight function  $w: E(G) \rightarrow \mathbb{R}$  and a walk  $W$  in  $G$ , let  $w(W)$  denote the sum of the weights along  $W$ , where the weight of an edge counts negatively when followed against its direction in  $\vec{G}$ .

A *tension* on  $\vec{G}$  is a weight function  $w$  such that  $w(C) = 0$  for every cycle  $C$  in  $G$ . Given a real number  $r$  with  $r \geq 2$ , an  $r$ -*tension* is a tension  $w$  such that  $1 \leq |w(uv)| \leq r - 1$  for every  $uv \in E(G)$ . An  $r$ -coloring  $f$  of  $G$  generates an  $r$ -tension  $w$  on an orientation  $\vec{G}$  by letting  $w(uv) = f(v) - f(u)$  for each  $uv \in E(\vec{G})$ .

A *modular  $r$ -tension* on an orientation  $\vec{G}$  is a weight function  $w: E(G) \rightarrow \mathbb{R}$  such that (1)  $w(C)$  is a multiple of  $r$  whenever  $C$  is a cycle in  $G$ , and (2) the weight on each edge differs by at least 1 from any multiple of  $r$ . Every  $r$ -tension is a modular  $r$ -tension, so an  $r$ -coloring of  $G$  generates a modular  $r$ -tension on  $\vec{G}$  as above.

Conversely, a modular  $r$ -tension  $w$  on  $\vec{G}$  generates an  $r$ -coloring  $f$  of  $G$  as follows. We may assume that  $G$  is connected (else do this in each component). Fix a vertex  $x$ . For each vertex  $v$ , choose an  $x, v$ -walk  $W$  in  $G$ , and choose  $f(v) \equiv w(W) \bmod r$  with  $0 \leq f(v) < r$ .

Since  $w$  is a modular  $r$ -tension,  $f(v)$  does not depend on the choice of  $W$ , and the colors on adjacent vertices differ by at least 1. We call the resulting  $f$  an  $r$ -coloring *generated from*  $w$ . We say “an” here because the coloring depends on the choice of  $x$ , but only by a cyclic permutation. We have shown that  $\chi_c(G)$  equals the least  $r$  such that some orientation  $\vec{G}$  has a modular  $r$ -tension.

Our lower bound on  $\chi'_c(H \square C_{2m+1})$  uses an analogue of girth, employing a parameter obtained from the cycle space of the graph. We obtain a strong lower bound when all the cycles in some basis of the cycle space are short.

Within the binary vector space of dimension  $|E(G)|$  with canonical basis vectors indexed by the edges, the *cycle space* of an undirected graph  $G$  is the subspace spanned by the incidence vectors of the cycles. The analogue for an orientation  $\vec{G}$  is the real vector space spanned by the signed incidence vectors of the cycles. For each cycle  $C$  in  $G$ , followed in a given direction, the *signed incidence vector* relative to  $\vec{G}$  has 1 or  $-1$  in each position for an edge of  $C$ , using  $-1$  if and only if the edge is followed against its direction in  $\vec{G}$ .

For any orientation  $\vec{G}$ , the same sets of cycles form bases of its cycle space as form bases of the cycle space of the underlying graph  $G$ . In either context, the number of nonzero positions in the incidence vector for a cycle is the same. Hence we define the relevant parameter in terms of  $G$ . For a basis  $\mathcal{B}$  of the cycle space of  $G$ , let  $\lambda(\mathcal{B})$  denote the maximum length of an element of  $\mathcal{B}$ . Let  $\lambda(G)$  denote the minimum of  $\lambda(\mathcal{B})$  over all bases of the cycle space. Note that  $\lambda(G)$  may be larger than the girth of  $G$ , but never smaller. The smaller the value of  $\lambda(G)$ , the larger the lower bound we will obtain on  $\chi_c(G)$ .

Before embarking on the technical lemmas, we pause to motivate their hypotheses. Let  $F = H \square C_{2m+1}$ . When  $H$  is  $(s-2)$ -regular,  $F$  is  $s$ -regular. Furthermore, the edges incident to any vertex of  $F$  become an  $s$ -clique in  $L(F)$ . Conversely, any two adjacent vertices of  $L(F)$  correspond to two incident edges in  $F$  and hence lie in an  $s$ -clique in  $L(F)$ . Therefore, we can study  $r$ -edge-colorings of  $F$  by studying  $r$ -colorings of  $L(F)$ , which we do by studying  $r$ -colorings of graphs in which every edge lies in a complete subgraph of order  $s$ .

**Lemma 2.2** *Let  $G$  be a graph such that each edge lies in a complete subgraph of order  $s$ . Let  $G$  have an  $r$ -coloring  $f$  such that*

$$r < s + \frac{1}{\lfloor \lambda(G)(1 - 1/s) \rfloor}.$$

*If  $\vec{G}$  is an orientation of  $G$ , then setting  $w(xy) = \lfloor \ell[f(x), f(y)]_r \rfloor$  for all  $xy \in E(\vec{G})$  defines a modular  $s$ -tension on  $\vec{G}$ .*

**Proof.** Let  $\epsilon = r - s$ , so  $\epsilon < \lfloor \lambda(G)(1 - 1/s) \rfloor^{-1}$ . For an edge  $xy$ , let  $Q$  be an  $s$ -clique containing  $x$  and  $y$ , and let  $f(Q) = \{f(v) : v \in Q\}$ . Let  $t = |f(Q) \cap [f(x), f(y)]_r|$ , so  $s - t = |f(Q) \cap [f(y), f(x)]_r|$ . By Lemma 2.1,

$$t \leq \ell([f(x), f(y)]_r) \leq t + \epsilon \quad \text{and} \quad s - t \leq \ell([f(y), f(x)]_r) \leq s - t + \epsilon. \quad (1)$$

By definition,  $w(xy) = \lfloor \ell([f(x), f(y)]_r) \rfloor$ , so

$$f(y) \in [f(x) + w(xy), f(x) + w(xy) + \epsilon]_r. \quad (2)$$

By (1),  $w(yx) = s - w(xy)$ . Let  $\mathcal{B}$  be a basis of the cycle space such that  $\lambda(\mathcal{B}) = \lambda(G) = k$ . To prove that  $w$  is a modular  $s$ -tension (when restricted to an orientation  $\vec{G}$  of  $G$ ), it suffices to show that  $w(C) \equiv 0 \pmod{s}$  for each  $C$  whose signed incidence vector lies in  $\mathcal{B}$ . (Since  $w(yx) = s - w(xy)$ , the choice of  $\vec{G}$  does not matter.)

Let  $x_0, \dots, x_{l-1}$  be the vertices of  $C$  in order, and let  $x_l = x_0$ ; note that  $l \leq k$ . Let  $e_i = x_i x_{i+1}$ . In testing whether  $w(C) \equiv 0 \pmod{s}$ , the orientation of the edges along  $C$  does not matter; all orientations yield the same congruence class for  $w(C)$ . Since the same sets of cycles yield bases under each orientation, in studying  $C$  we may assume an orientation with each  $e_i$  directed from  $x_i$  to  $x_{i+1}$ . Now  $w(C) \equiv \sum_{i=0}^{l-1} w(e_i) \pmod{s}$ .

Since each edge lies in a complete subgraph of order  $s$ , (2) applies to each edge, so  $f(x_{i+1}) \in [f(x_i) + w(e_i), f(x_i) + w(e_i) + \epsilon]_r$  for  $0 \leq i \leq l-1$ . Combining the allowed variations in the intervals for all edges of  $C$  yields

$$f(x_0) \in [f(x_0) + w(C), f(x_0) + w(C) + l\epsilon]_r. \quad (3)$$

By symmetry, we may choose  $f(x_0) = 0$ , which reduces (3) to  $0 \in [w(C), w(C) + l\epsilon]_r$ .

Since  $1 \leq w(e_i) \leq s-1$ , we have  $l \leq w(C) \leq (s-1)l$ . Since  $w(C)$  is an integer, by choosing  $q$  to be  $\lceil w(C)/s \rceil$  or  $\lfloor w(C)/s \rfloor$  we can write  $w(C) = qs + j$  for integers  $q$  and  $j$  such that  $\lceil l/s \rceil \leq q \leq \lfloor l(1-1/s) \rfloor$  and  $|j| \leq s-1$ . Now

$$[w(C), w(C) + l\epsilon]_r = [j + qr - q\epsilon, j + qr + (l-q)\epsilon]_r = [j - q\epsilon, j + (l-q)\epsilon]_r.$$

Since  $q \leq \lfloor l(1-1/s) \rfloor$ , we have  $q\epsilon \leq \lfloor l(1-1/s) \rfloor \epsilon \leq \lfloor k(1-1/s) \rfloor \epsilon < 1$ . Similarly,  $q \geq \lceil l/s \rceil$  yields  $(l-q)\epsilon \leq \lfloor l(1-1/s) \rfloor \epsilon < 1$ . Since  $0 \in [w(C), w(C) + l\epsilon]_r \subseteq (j-1, j+1)_r$ , we thus have  $j = 0$ . That is,  $w(C) \equiv 0 \pmod{s}$ . Thus  $w$  is a modular  $s$ -tension on  $\vec{G}$ .  $\blacksquare$

The conclusion of Lemma 2.2 states that  $G$  is  $s$ -colorable. This is impossible if  $G$  is  $s$ -regular with odd order, so the lemma implies that  $\chi_c(G) \geq s + \lfloor \lambda(G)(1-1/s) \rfloor^{-1}$ . With  $G = L(H \square C_{2m+1})$ , we obtain a lower bound for  $\chi'_c(H \square C_{2m+1})$ , but it is not the lower bound we seek. The cycle space for  $G$  contains copies of the cycle space for  $H$ , but it is larger, and it may be that  $\lambda(G) > \lambda(H)$ , so the bound may be weaker than desired. To improve the bound, we will study subgraphs of  $G$  where we can control the value of  $\lambda$ . Before introducing these subgraphs, we prove a technical lemma about the color classes of the colorings generated from the modular  $s$ -tension produced by Lemma 2.2.

**Lemma 2.3** *Let  $G$  be a graph such that each edge lies in a complete subgraph of order  $s$ . Suppose that  $G$  has an  $r$ -coloring  $f$  such that*

$$r < s + \frac{1}{\lfloor \lambda(G)(1-1/s) \rfloor}.$$

For a fixed vertex  $v^* \in V(G)$  and any  $x \in V(G)$ , let  $g(x) = \lfloor \ell([f(v^*), f(x)]_r) \rfloor$ . This function  $g$  is a proper (integer)  $s$ -coloring of  $G$  that satisfies the following property:  $g(x) = g(x')$  if and only if  $G$  has a vertex list  $(x, \dots, x')$  in which any consecutive entries  $v$  and  $v'$  satisfy  $d_G(v, v') = 2$  and  $|f(v) - f(v')|_r < 1/2$ .

**Proof.** Call a list  $(x, \dots, x')$  with the specified properties an  $x, x'$ -skiplist.

Let  $\vec{G}$  be an orientation of  $G$ . By Lemma 2.2, setting  $w(xy) = \lfloor \ell([f(x), f(y)]_r) \rfloor$  for all  $xy \in E(\vec{G})$  defines a modular  $s$ -tension  $w$  on  $\vec{G}$ , and  $g$  is an  $s$ -coloring of  $G$  generated from  $w$ . Since the values of  $w$  are integers in  $\{0, \dots, s-1\}$ , in fact  $g$  is a proper (integer)  $s$ -coloring of  $G$ .

Vertices  $y$  and  $y'$  with  $|f(y) - f(y')|_r < 1/2$  must be nonadjacent. If they have a common neighbor  $z$ , then

$$f(z) \in [f(y) + w(yz), f(y) + w(yz) + \epsilon]_r \cap [f(y') + w(y'z), f(y') + w(y'z) + \epsilon]_r.$$

If  $|w(y'z) - w(yz)| \geq 1$ , then the intervals on the right are disjoint, since  $\epsilon < 1/2$  and  $|f(y) - f(y')|_r < 1/2$ . Therefore  $w(yz) = w(y'z)$ , which yields  $g(y) = g(y')$ . Therefore, all vertices in an  $x, x'$ -skiplist have the same color under  $g$ ; in particular,  $g(x) = g(x')$ .

Conversely, suppose that  $g(x) = g(x')$ . Let  $v_0, \dots, v_t$  be the vertices along an  $x, x'$ -path in  $G$ , with  $x = v_0$  and  $x' = v_t$ . For  $0 \leq i \leq t-1$ , let  $X_i$  be an  $s$ -clique of  $G$  containing  $v_i$  and  $v_{i+1}$ . Select auxiliary vertices  $x_0, \dots, x_t$  as follows. Having selected  $x_0, \dots, x_{i-1}$  (starting with  $x_0 = v_0 = x$ ), observe that  $v_i \in X_{i-1} \cap X_i$ . By Lemma 2.1, there is a unique vertex  $x_i \in X_i$  with  $|f(x_i) - f(x_{i-1})|_r \leq \epsilon < 1/2$ . Applying the preceding paragraph with  $y = x_i$  and  $y' = x_{i-1}$  yields  $g(x_i) = g(x_{i-1})$ . Finally,  $x_t = x'$ , since  $x_t, x' \in X_t$  and  $g(x_t) = g(x) = g(x')$ . Now  $(x_0, \dots, x_t)$  is an  $x, x'$ -skiplist.  $\blacksquare$

The crucial consequence of Lemma 2.3 is that the partition of  $G$  into color classes under  $g$  does not depend on the choice of  $v^*$ .

### 3 A Lower Bound on $\chi'_c(H \square C_{2m+1})$

We specialize again to the study of  $\chi'_c(H \square C_{2m+1})$ . When  $H$  is  $(s-2)$ -regular with odd order, the product  $H \square C_{2m+1}$  is  $s$ -regular with odd order and hence is Class 2. Thus  $\chi'_c(H \square C_{2m+1}) > s$ . We improve this lower bound when  $s$  is divisible by 4.

Let  $V(C_{2m+1}) = \{v_0, \dots, v_{2m}\}$ , indexed in order; treat subscripts modulo  $2m+1$ . The  $i$ th layer  $H_i$  of  $H \square C_{2m+1}$  is the subgraph induced by  $V(H) \times \{v_i\}$ . Each layer  $H_i$  is isomorphic to  $H$ . For  $e \in E(H)$  and  $x \in V(H)$ , let  $e^i$  and  $x^i$  denote the copies of  $e$  and  $x$  in  $H_i$ . We call  $\bigcup_{i=0}^{2m} E(H_i)$  the *horizontal edges* of  $H \square C_{2m+1}$ .

For  $x \in V(H)$ , let  $l_x^i$  denote the edge  $x^i x^{i+1}$  in  $H \square C_{2m+1}$ . Let  $L_i = \{l_x^i : x \in V(H)\}$ ; we call  $L_i$  the  $i$ th *link* of  $H \square C_{2m+1}$  and call  $\bigcup_{i=0}^{2m} L_i$  the *vertical edges* of  $H \square C_{2m+1}$ .

In a graph  $G$  whose vertices all have degree  $s$  or  $1$ , any two incident edges are incident at a vertex of degree  $s$ . Therefore, in  $L(G)$  every edge lies in a complete subgraph of order  $s$ . We will be applying the results of Section 2 to subgraphs of  $H \square C_{2m+1}$  having the form  $L_{i-1} \cup H_i \cup L_i$ , where every vertex has degree  $s$  or  $1$ . We also need the following observation.

**Lemma 3.1** *For any graph  $G$ , the equality  $\lambda(L(G)) = \lambda(G)$  holds.*

**Proof.** Since cycles in  $G$  turn into cycles in  $L(G)$  and must be spanned by any basis for  $L(G)$ , we have  $\lambda(L(G)) \geq \lambda(G)$ . Also, a basis for the cycle space of  $G$  (indexed by edges) can be augmented to a basis for the cycle space of  $L(G)$  (indexed by vertices) by adding the incidence vectors of triangles in  $L(G)$  consisting of three edges in  $G$  having a common endpoint. The added vectors have weight 3, so  $\lambda(L(G)) \leq \lambda(G)$ . ■

**Theorem 3.2** *If  $H$  is an  $(s-2)$ -regular graph of odd order, where  $4 \mid s$ , then*

$$\chi'_c(H \square C_{2m+1}) \geq s + \frac{1}{\lfloor \lambda(H)(1 - 1/s) \rfloor}.$$

**Proof.** If not, then  $H \square C_{2m+1}$  has an  $(s + \epsilon)$ -edge-coloring  $f$ , where  $\epsilon < \lfloor \lambda(H)(1 - 1/s) \rfloor^{-1}$ .

Let  $G_i$  be the subgraph of  $L(H \square C_{2m+1})$  induced by  $L_{i-1} \cup E(H_i) \cup L_i$  (as defined above). Each edge of  $G_i$  lies in a complete subgraph of order  $s$ . Let  $\mathcal{T}$  be the set of triangles in  $G_i$ . If  $\mathcal{B}$  is a basis of the cycle space of  $L(H_i)$ , then  $\mathcal{B} \cup \mathcal{T}$  contains a basis of the cycle space of  $G_i$ . Thus  $\lambda(G_i) = \lambda(L(H_i)) = \lambda(L(H)) = \lambda(H)$ , using  $H_i \cong H$  and Lemma 3.1.

For each  $G_i$ , Lemma 2.3 states that the function  $g_i$  defined by fixing  $v^* \in V(G_i)$  and setting  $g_i(x) = \lfloor \ell([f(v^*), f(x)]_r) \rfloor$  for all  $x \in V(G_i)$  is a proper (integer)  $s$ -coloring of  $G_i$ . Since this  $g_i$  depends only on the global  $r$ -coloring  $f$  and the choice of  $v^*$ , the restrictions to  $L_i$  of the partitions of  $V(G_i)$  and  $V(G_{i+1})$  into color classes under  $g_i$  and  $g_{i+1}$  are the same when  $v^*$  is chosen to be an element of  $L_i$ .

Furthermore, Lemma 2.3 implies that the partition of  $V(G_i)$  into color classes does not depend on the choice of  $v^*$ ; it is determined only by values of  $f$  and distances between vertices in  $G_i$ . We conclude that no matter how  $v_i^*$  and  $v_{i+1}^*$  are chosen in specifying  $g_i$  and  $g_{i+1}$ , the resulting partitions of  $L_i$  into color classes are the same.

Each vertex  $x^i$  of the product has two incident vertical edges, namely  $l_x^i$  and  $l_x^{i-1}$ . We say that a color  $j$  is a *vertical color at  $x^i$*  if some vertical edge incident to  $x^i$  has color  $j$  under  $g_i$ . For each  $x^i \in V(H_i)$ , the  $s$  incident edges of  $G_i$  have distinct colors. Therefore a color  $j$  is a vertical color at  $x^i$  if and only if no edge of  $H_i$  incident to  $x^i$  has color  $j$  under  $g_i$ . Since  $H$  has odd order, and the number of vertices of  $H_i$  incident to edges of  $H_i$  with color  $j$  is even, we conclude that  $j$  is a vertical color at an odd number of vertices of  $H_i$ . In other words, in the partition of  $L_{i-1} \cup L_i$  formed by the color classes under  $g_i$ , each class has odd size.

Let  $C_i^+$  [respectively,  $C_i^-$ ] be the set of colors used by  $g_i$  on an odd number of edges of  $L_i$  [respectively,  $L_{i-1}$ ]. Since each class under  $g_i$  has odd size in  $L_i \cup L_{i-1}$ , we conclude that  $j \in C_i^-$  if and only if  $j \notin C_i^+$ .

Since  $|L_i|$  and  $|L_{i-1}|$  are odd, it follows that  $|C_i^+|$  and  $|C_i^-|$  are also odd. Since  $|C_i^+| + |C_i^-| = s$  and  $s$  is divisible by 4, it follows that  $|C_i^+| \neq |C_i^-|$ . Since  $g_i$  and  $g_{i+1}$  induce the same partitions of  $L_i$ , it follows that  $|C_{i+1}^-| = |C_i^+|$ , and hence also  $|C_{i+1}^+| = |C_i^-|$ . Now the values of  $|C_i^+|$  must alternate between two distinct values as  $i$  runs through all  $2m + 1$  subscripts, which is impossible since  $2m + 1$  is odd.  $\blacksquare$

## 4 An Upper Bound on $\chi'_c(H \square C_{2m+1})$

In this section, we obtain an upper bound on  $\chi'_c(H \square C_{2m+1})$  for some  $H$ . As a consequence, we show that  $\chi'_c(H \square C_{2m+1}) - \Delta(H \square C_{2m+1})$  can be bounded above by a number that is arbitrarily close to  $\chi'_c(H) - \Delta(H)$  by making  $m$  sufficiently large.

We show first that increasing  $m$  cannot increase the circular chromatic index. We simply use the coloring of one layer on three consecutive layers in the larger graph and re-use the colorings on its neighboring links.

**Lemma 4.1** *If  $m' \geq m$ , then  $\chi'_c(H \square C_{2m'+1}) \leq \chi'_c(H \square C_{2m+1})$ .*

**Proof.** It suffices to prove that  $\chi'_c(H \square C_{h+2}) \leq \chi'_c(H \square C_h)$  for all  $h$ . Let  $f$  be an  $r$ -edge-coloring of  $H \square C_h$ . Form an  $r$ -edge-coloring of  $H \square C_h$  as follows. Color the layers  $H_0, \dots, H_{h-1}$  and links  $L_0, \dots, L_{h-1}$  as under  $f$ . Color the layers  $H_h$  and  $H_{h+1}$  the same as  $H_{h-1}$ . Color the links  $L_h$  and  $L_{h+1}$  the same as  $L_{h-2}$  and  $L_{h-1}$ , respectively. Now the colors on any two incident edges of  $H \square C_{h+2}$  under  $f'$  are also colors on two incident edges of  $L(H \square C_h)$  under  $f$ . Thus  $f'$  is also an  $r$ -edge-coloring.

The colors on any two adjacent vertices of  $L(H \square C_{2m+3})$  under  $f'$  are also colors on two adjacent vertices of  $L(H \square C_{2m+1})$  under  $f$ . Thus  $f'$  is also an  $r$ -coloring.  $\blacksquare$

Since  $\chi'_c(H \square C_{2m+1}) \geq \Delta(H \square C_{2m+1}) = \Delta(H) + 2$  for all  $m$ , Lemma 4.1 implies that  $\chi'_c(H \square C_{2m+1})$  has a limit as  $m \rightarrow \infty$ . In Section 5 we show that this limit is attained when  $H$  is an odd cycle, and we compute its value.

To prove the upper bound, we need a standard result about circular coloring.

**Lemma 4.2** (See [13]) *If a graph  $G$  has a  $r$ -coloring  $f$  with  $r = p/q$  where  $p, q \in \mathbb{N}$ , then it has an  $r$ -coloring  $f'$  such that the colors under  $f'$  are multiples of  $1/q$ , and such that if  $xy \in E(G)$ , then  $|f'(x) - f'(y)|_r$  differs by less than  $1/q$  from  $|f(x) - f(y)|_r$ .*

**Proof.** Let  $f'(x) = \lfloor qf(x) \rfloor / q$  (such multiplication arguments were used as early as [3]). Note that  $f'(x)$  is the largest multiple of  $1/q$  that does not exceed  $f(x)$ . Under this transformation,  $|f'(x) - f'(y)|_r$  equals  $|f(x) - f(y)|$  if the latter is a multiple of  $1/q$ . Otherwise, the difference shifts to the next larger or next smaller multiple of  $1/q$ .

In particular, if the colors assigned to two vertices differ by at least  $a/q$  before the transformation, for some positive integer  $a$  (such as  $a = q$ ), then they also differ by at least  $a/q$  after the transformation. Thus  $f'$  is an  $r$ -coloring.  $\blacksquare$

Given an  $r$ -edge-coloring of a graph  $H$ , a *color gap* for a vertex  $x$  of  $H$  is a maximal open interval on the circle  $C^r$  that contains no color used on an edge incident to  $x$ .

**Theorem 4.3** *Let  $H$  be a graph having a  $p/q$ -edge-coloring  $f$  such that every vertex  $x$  of  $H$  has a color gap of length at least 3. If  $p$  is odd and  $2m + 1 \geq p$ , then  $\chi'_c(H \square C_{2m+1}) \leq p/q$ .*

**Proof.** By Lemma 4.1, it suffices to prove this when  $2m + 1 = p$ . By Lemma 4.2 (applied to  $L(H)$ ), we may assume that each  $f(e)$  is a multiple of  $1/q$ , still with each vertex having a color gap of length at least 3 (using  $a = 3q$  in that argument). For each  $x \in V(H)$ , let  $(a_x, b_x)_{p/q}$  be a color gap under  $f$  with length at least 3.

We produce a  $p/q$ -edge-coloring  $\phi$  for  $H \square C_{2m+1}$ . We use the same coloring  $f$  in each layer, except that the colors in each layer increase by one unit from the colors on the corresponding edges in the previous layer. Since  $2m = p - 1 = q(p/q) - 1$ , the colors on layer  $H_0$  are also one unit (modulo  $p/q$ ) above the corresponding colors on  $H_{2m}$ . This is achieved by letting  $\phi(e^i) = f(e) + i \pmod{p/q}$  for each  $e \in E(H)$  and  $0 \leq i \leq 2m$ .

It now suffices to use the color gaps to fit in colors for the vertical edges. Specifically, we set  $\phi(l_x^i) = a_x + 2 + i \pmod{p/q}$  for each  $x \in V(H)$  and  $0 \leq i \leq 2m$ . Since no horizontal edge at  $x^i$  receives a color in  $(a_x + i, a_x + i + 3)$ , the colors  $a_x + i + 1$  and  $a_x + i + 2$  are available for  $l_x^{i-1}$  and  $l_x^i$ , respectively, when viewed from  $x^i$ . Furthermore,  $\phi$  achieves this assignment simultaneously for the vertical edges at all  $x^j$ . Hence for all incident edges, the assigned colors differ by at least 1.  $\blacksquare$

For any graph  $G$ , let  $\partial(G) = \chi'_c(G) - \Delta(G)$ . Thus  $G$  is Class 1 if and only if  $\partial(G) = 0$ , and otherwise  $0 < \partial(G) \leq 1$ .

**Corollary 4.4** *For any graph  $H$ ,  $\lim_{m \rightarrow \infty} \partial(H \square C_{2m+1}) \leq \partial(H)$ .*

**Proof.** The limit exists, using  $\Delta(H \square C_{2m+1}) \geq \Delta(H) + 2$  and Lemma 4.1. It suffices to show, given  $\epsilon > 0$ , that  $\partial(H \square C_{2m+1}) \leq \partial(H) + \epsilon$  when  $m$  is sufficiently large.

Choose  $p, q \in \mathbb{N}$  with  $p$  odd such that  $\chi'_c(H) \leq p/q \leq \chi'_c(H) + \epsilon$ . Let  $f$  be a  $p/q$ -edge-coloring of  $H$ . Also  $f$  can be viewed as a  $(p/q + 2)$ -edge-coloring of  $H$ . For  $x \in V(H)$ , let  $b_x$  and  $a_x$  be the minimum and maximum colors in  $[0, p/q)$  used on edges incident to  $x$ , respectively. Since  $\ell((a_x, b_x)_{p/q}) \geq 1$ , also  $\ell((a_x, b_x)_{p/q+2}) \geq 3$ . Relative to  $f$  as a  $(p/q + 2)$ -edge-coloring, each vertex of  $H$  thus has a color gap of length at least 3. By Theorem 4.3,  $\chi'_c(H \square C_{2m+1}) \leq p/q + 2 \leq \Delta(H \square C_{2m+1}) + \epsilon$  when  $2m + 1 \geq p$ .  $\blacksquare$

Recall that  $H \square H'$  is Class 1 when  $H$  or  $H'$  is Class 1. That is,  $\partial(H) = 0$  or  $\partial(H') = 0$  implies  $\partial(H \square H') = 0$ . It is natural to ask if  $\partial(H \square H') \leq \min\{\partial(H), \partial(H')\}$  always holds.

It does not, by the following example. Let  $H = C_{2k+1}$  and  $H' = C_{2m+1}$ . Since  $\chi'_c(C_{2m+1}) = 2 + 1/m$ , we can make  $\partial(H')$  arbitrarily small. However,  $\lambda(H) = 2k + 1$ , so Theorem 3.2 yields  $\partial(H \square H') \geq \lfloor (6k + 3)/4 \rfloor^{-1} = \lceil 3k/2 \rceil^{-1}$ , independent of  $m$ .

On the other hand,  $\lceil 3k/2 \rceil^{-1} < k^{-1} = \partial(C_{2k+1})$ . Based on this and Theorem 4.3 and other examples, we propose the following conjecture.

**Conjecture 4.5** *For any graphs  $H$  and  $H'$ ,  $\partial(H \square H') \leq \max\{\partial(H), \partial(H')\}$ .*

## 5 Tightness of the lower bound

As noted above, Theorem 3.2 implies that  $\chi'_c(C_{2k+1} \square C_{2m+1}) \geq 4 + \lceil 3k/2 \rceil^{-1}$  for all  $m$ . In this section, we prove that the bound is sharp when  $m \geq 3k + 1$ . This proves Conjecture 4.5 for products of two odd cycles when one is at least three times as long as the other.

**Lemma 5.1** *If there exist integers  $\alpha, \beta, q$  with  $0 < q \leq m/2$  such that  $|\alpha| + |\beta| = 2k + 1$  and  $\alpha q + \beta(q + 1) \equiv 0 \pmod{4q + 1}$ , then  $\chi'_c(C_{2k+1} \square C_{2m+1}) \leq 4 + 1/q$ .*

**Proof.** By Theorem 4.3 with  $p = 4q + 1$ , it suffices to produce a  $(4 + 1/q)$ -edge-coloring  $f$  of  $C_{2k+1}$  such that every vertex  $x$  of  $C_{2k+1}$  has a color gap of length at least 3. Since  $C_{2k+1}$  is 2-regular, and we use a color circle of length  $4 + 1/q$ , the condition on  $f$  becomes “If  $e$  and  $e'$  are incident edges in  $C_{2k+1}$ , then  $1 \leq |f(e') - f(e)|_{(4+1/q)} \leq 1 + 1/q$ .” Multiplying by  $q$ , we further transform this to seeking integers  $z_1, \dots, z_{2k+1}$  modulo  $4q + 1$  such that neighboring integers differ by  $q$  or  $q + 1$ .

In the hypothesis, we may assume by symmetry that  $\alpha \geq 0$ . We construct the first  $\alpha$  and last  $|\beta|$  integers as separate arithmetic progressions, with common difference  $q$  for the first  $\alpha$  and  $q + 1$  for the last  $|\beta|$ . For  $1 \leq i \leq \alpha$ , let  $z_i = iq$  (this portion is empty if  $\alpha = 0$ ). For  $1 \leq i \leq |\beta|$ , let  $z_{\alpha+i} = \alpha q + \epsilon i(q + 1)$ , where  $\epsilon = 1$  if  $\beta > 0$  and  $\epsilon = -1$  if  $\beta < 0$ .

The construction enforces the needed differences until just before the end; we need only compare  $z_{2k+1}$  and  $z_1$ . Since  $z_{2k+1} = \alpha q + \beta(q + 1) \equiv 0 \pmod{4q + 1}$ , indeed  $z_{2k+1}$  and  $z_1$  differ by  $q$ . ■

**Theorem 5.2** *If  $m \geq 3k + 1$ , then  $\chi'_c(C_{2k+1} \square C_{2m+1}) = 4 + \lceil 3k/2 \rceil^{-1}$ .*

**Proof.** We have noted that Theorem 3.2 gives the lower bound. It suffices to find integers  $\alpha, \beta, q$  satisfying the hypotheses of Lemma 5.1 with  $q = \lceil 3k/2 \rceil = \lfloor (6k + 3)/4 \rfloor$ .

Let  $r = \lfloor (k - 1)/2 \rfloor$ , so  $k = 2r + s$  with  $1 \leq s \leq 2$ . Now  $q = 3r + s + 1$ . Let  $\alpha = s - 1$  and  $\beta = -(4r + s + 2)$ . We have  $|\alpha| + |\beta| = (4r + 2s + 1) = 2k + 1$  and

$$\alpha q + \beta(q + 1) = (s - 1)q - (4r + s + 2)(q + 1) = -(4q + 1)(r + 1),$$

where the last computation uses  $q = 3r + s + 1$ . Thus  $\alpha q + \beta(q + 1) \equiv 0 \pmod{4q + 1}$ , and Lemma 5.1 applies. ■

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