

# Coloring the Square of a $K_4$ -minor Free Graph

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Revised, December 24, 2002

## Abstract

Let  $G$  be a  $K_4$ -minor free graph with maximum degree  $\Delta$ . We prove that the chromatic number of the square of  $G$  is at most (i)  $\Delta + 3$  if  $2 \leq \Delta \leq 3$ ; or (ii)  $\lfloor 3\Delta/2 \rfloor + 1$  if  $\Delta \geq 4$ . Examples are given to show the bounds can be attained.

*Key words:* Series-parallel graph,  $K_4$ -minor free graph, square, chromatic number

## 1 Introduction

All graphs considered in this paper are finite, loopless, and without multiple edges. For a graph  $G$ , let  $V(G)$ ,  $E(G)$ ,  $|G|$ ,  $\Delta(G)$ , and  $\delta(G)$  denote, respectively, its vertex set, edge set, number of vertices, maximum degree, and

minimum degree. For  $x \in V(G)$ , let  $N_G(x)$  denote the set of neighbors of  $x$  in  $G$  and let  $d_G(x)$  denote the degree, i.e., the number of neighbors, of  $x$  in  $G$ . A vertex of degree  $k$  is called a  $k$ -vertex. For two vertices  $u, v \in V(G)$ , let  $\text{dist}_G(u, v)$  denote the distance between  $u$  and  $v$ , i.e., the length of a shortest path connecting them.

A  $k$ -coloring of a graph  $G$  is a mapping  $\phi$  from  $V(G)$  to the set of colors  $\{1, 2, \dots, k\}$  such that  $\phi(x) \neq \phi(y)$  for every edge  $xy$  of  $G$ . We call  $G$   $k$ -colorable if it has a  $k$ -coloring. The *chromatic number*  $\chi(G)$  of  $G$  is the smallest integer  $k$  such that  $G$  is  $k$ -colorable. The *square*  $G^2$  of a graph  $G$  is the graph defined on the vertex set  $V(G)$  such that two vertices  $u$  and  $v$  are adjacent in  $G^2$  if and only if  $1 \leq \text{dist}_G(u, v) \leq 2$ . A mapping  $\phi$  is a  $k$ -coloring of  $G^2$  if and only if  $\phi(u) \neq \phi(v)$  whenever  $1 \leq \text{dist}_G(u, v) \leq 2$ .

Since  $G^2$  contains a clique of order at least  $\Delta(G) + 1$ , we have  $\chi(G^2) \geq \Delta(G) + 1$  for any graph  $G$ . If  $G$  is a tree, then  $\chi(G^2) = \Delta(G) + 1$ . On the other hand,  $\Delta(G^2) \leq \Delta^2(G)$ . Hence  $\chi(G^2) \leq \Delta^2(G) + 1$  for any graph  $G$ . The 5-cycle  $C_5$  and the Petersen graph satisfy  $\chi(G^2) = \Delta^2(G) + 1$ .

Wegner [9] first investigated the chromatic number of the square of a planar graph. Wegner proved that  $\chi(G^2) \leq 8$  for every planar graph  $G$  with  $\Delta(G) = 3$  and conjectured that the upper bound could be reduced to 7, which has been confirmed by Thomassen [7]. For planar graphs of maximum degree  $\Delta \geq 4$ , Wegner [9] proposed the following conjecture. The bounds are sharp if the conjecture is true.

**Conjecture 1** *Let  $G$  be a planar graph. Then*

$$\chi(G^2) \leq \begin{cases} \Delta(G) + 5 & \text{if } 4 \leq \Delta(G) \leq 7; \\ \lfloor 3\Delta(G)/2 \rfloor + 1 & \text{if } \Delta(G) \geq 8. \end{cases}$$

This conjecture remains open. The best upper bound so far is  $5\Delta(G)/3 + 78$  established by Molloy and Salavatipour [5]. This improves other recently

obtained upper bounds:  $\lfloor 9\Delta/5 \rfloor + 2$  for  $\Delta(G) \geq 749$  ([1]),  $\lfloor 9\Delta/5 \rfloor + 1$  for  $\Delta(G) \geq 47$  ([2]), and  $2\Delta(G) + 25$  ([6]). For planar graphs of large girth, better upper bounds for  $\chi(G^2)$  are known. Wang and Lih [8] proved that if  $G$  is a planar graph with girth  $g(G)$ , then  $\chi(G^2) \leq \Delta(G) + 5$  when  $g(G) \geq 7$ ,  $\chi(G^2) \leq \Delta(G) + 10$  when  $g(G) \geq 6$ , and  $\chi(G^2) \leq \Delta(G) + 16$  when  $g(G) \geq 5$ .

A graph  $G$  has a graph  $H$  as a *minor* if  $H$  can be obtained from a subgraph of  $G$  by contracting edges, and  $G$  is called *H-minor free* if  $G$  does not have  $H$  as a minor. A graph  $G$  is called a *series-parallel* graph if  $G$  can be obtained from  $K_2$  by applying a sequence of operations, where each operation is either to duplicate an edge (i.e., replace an edge with two parallel edges) or to subdivide an edge (i.e., replace an edge with a path of length 2). It is well-known [3] that a graph  $G$  is an outerplanar graph if and only if  $G$  is  $K_4$ -minor free and  $K_{2,3}$ -minor free. A graph  $G$  is  $K_4$ -minor free if and only if each block of  $G$  is a series-parallel graph. Thus the class of  $K_4$ -minor free graphs is a class of planar graphs that contains both outerplanar graphs and series-parallel graphs. We will establish the following in this paper.

**Theorem 1** *Let  $G$  be a  $K_4$ -minor free graph. Then*

$$\chi(G^2) \leq \begin{cases} \Delta(G) + 3 & \text{if } 2 \leq \Delta(G) \leq 3; \\ \lfloor 3\Delta(G)/2 \rfloor + 1 & \text{if } \Delta(G) \geq 4. \end{cases}$$

Before proving Theorem 1, we show that the result is best possible. For  $\Delta = 2$ , we have  $\Delta(C_5) = 2$  and  $\chi(C_5^2) = 5$ . For  $\Delta = 3$ , let  $G$  be the graph consisting of three internally disjoint paths joining two vertices  $x$  and  $y$ , where two of the paths are of length 2, and the third one is of length 3. Then  $\Delta(G) = 3$  and  $\chi(G^2) = \chi(K_6) = 6$ . For  $\Delta = 2k \geq 4$ , let  $G_{2k}$  be the graph consisting of  $k$  internally disjoint paths joining  $x$  and  $y$ ,  $k$  internally disjoint paths joining  $x$  and  $z$ , and  $k$  internally disjoint paths joining  $y$  and  $z$ . All these paths are of length 2, except one path joining  $x$  and  $y$  is of

length 1, and one path joining  $x$  and  $z$  is of length 1. Then  $\Delta(G_{2k}) = 2k$  and  $\chi(G_{2k}^2) = \chi(K_{3k+1}) = 3k + 1$ . For  $\Delta = 2k + 1 \geq 5$ , let  $G_{2k+1}$  be obtained from  $G_{2k}$  by adding a new path of length 2 joining  $y$  and  $z$ . Then  $\Delta(G_{2k+1}) = 2k + 1$  and  $\chi(G_{2k+1}^2) = \chi(K_{3k+2}) = 3k + 2$ .

## 2 Proof of Theorem 1

Let  $u \sim v$  denote that  $u$  and  $v$  are adjacent in  $G$ , that is,  $uv \in E(G)$ . Define  $S_G(u) = \{x \mid d_G(x) \geq 3 \text{ such that } u \sim x \text{ or there exists a 2-vertex } z \text{ satisfying } u \sim z \text{ and } z \sim x\}$ . Let  $D_G(u) = |S_G(u)|$ . It is well-known [4] that every  $K_4$ -minor free graph contains a vertex of degree at most two.

**Lemma 2** *Let  $G$  be a  $K_4$ -minor free graph. Then one of the following holds.*

- (i)  $\delta(G) \leq 1$ ;
- (ii) *There exist two adjacent 2-vertices;*
- (iii) *There exists a vertex  $u$  with  $d_G(u) \geq 3$  such that  $D_G(u) \leq 2$ .*

**Proof.** Suppose that the lemma is false. Let  $G$  be a counterexample, i.e.,  $G$  is a  $K_4$ -minor free graph satisfying the following properties:

- (a)  $\delta(G) = 2$ ;
- (b) Every 2-vertex is adjacent to two vertices of degree  $\geq 3$ ;
- (c) Every vertex  $u$  with  $d_G(u) \geq 3$  has  $D_G(u) \geq 3$ .

Let  $H$  be the graph obtained from  $G$  by removing all 2-vertices and joining any two nonadjacent vertices that had a 2-vertex as a common neighbor. Then  $H$  has minimum degree at least 3 because  $N_H(v) = S_G(v)$  for each vertex  $v$  with  $d_G(v) \geq 3$ . However,  $H$  is a  $K_4$ -minor free graph by its very construction. A contradiction is henceforth obtained.  $\square$

**Proof of Theorem 1.** The case for  $\Delta(G) = 2$  follows directly from the well-known Brooks' Theorem since  $\Delta(G) = 2$  implies  $\Delta(G^2) \leq 4$ . In the

following we assume that  $\Delta(G) \geq 3$ .

Our proof proceeds by induction on the number of vertices of  $G$ . To make notation simpler, define  $K(\Delta) = 6$  if  $\Delta = 3$  and  $K(\Delta) = \lfloor 3\Delta/2 \rfloor + 1$  if  $\Delta \geq 4$ .

The conclusion follows immediately if  $|G| \leq 6$ . Now assume that  $|G| \geq 7$  and Theorem 1 holds for  $K_4$ -minor free graphs  $H$  with  $|H| < |G|$ .

If  $G$  has a vertex  $x$  of degree 1, then let  $H = G - x$ . If  $G$  has two adjacent 2-vertices  $x$  and  $y$ , then let  $H$  be the graph obtained from  $G$  by deleting  $x$  and adding an edge joining the two neighbors of  $x$ . In each of the above cases,  $H$  is a  $K_4$ -minor free graph with  $\Delta(H) \leq \Delta(G)$ . Moreover,  $G^2 - x$  is a subgraph of  $H^2$ . By the induction hypothesis,  $H^2$  is  $K(\Delta)$ -colorable. So  $G^2 - x$  is  $K(\Delta)$ -colorable. Since  $x$  has degree  $< K(\Delta)$  in  $G^2$ , any  $K(\Delta)$ -coloring of  $G^2 - x$  can be extended to a  $K(\Delta)$ -coloring of  $G^2$ .

Thus we may assume that  $\delta(G) = 2$  and any two 2-vertices of  $G$  are not adjacent. By Lemma 2,  $G$  has a vertex  $u$  such that  $d_G(u) \geq 3$  and  $D_G(u) \leq 2$ . For  $t \in S_G(u)$ , let  $M(u, t)$  denote the set of all 2-vertices in  $G$  that are adjacent to both  $u$  and  $t$  and let  $m(t) = |M(u, t)|$ .

Obviously  $D_G(u) \geq 1$ . Assume that  $D_G(u) = 1$  and  $S_G(u) = \{z\}$ . Then all the neighbors of  $u$  are either  $z$  or some neighbors of  $z$ . Since  $d_G(u) \geq 3$ , we see that  $m(z) \geq 2$ . Let  $w \in M(u, z)$  and  $H = G - w$ . Then  $H$  is a  $K_4$ -minor free graph with  $\Delta(H) \leq \Delta(G)$ . Moreover,  $G^2 - w = H^2$ . By the induction hypothesis,  $H^2$  has a  $K(\Delta)$ -coloring  $\phi$ . Since  $d_{G^2}(w) \leq \Delta(G) + 1 < K(\Delta)$ ,  $\phi$  can be extended to a  $K(\Delta)$ -coloring of  $G^2$ .

Assume that  $D_G(u) = 2$ . Let  $S_G(u) = \{x, y\}$ . Thus all the neighbors of  $u$  are either  $x, y$ , or some neighbors of  $x$  or  $y$ . Without loss of generality, we suppose  $m(x) \geq m(y)$ . Since  $d_G(u) \geq 3$ , we have  $m(x) \geq 1$ . Let  $w \in M(u, x)$ . The proof is divided into the following three cases.

**Case 1.**  $x \sim u$ .

Let  $H = G - w$ . Then  $H$  is a  $K_4$ -minor free graph with  $\Delta(H) \leq \Delta(G)$  and  $G^2 - w = H^2$ . By the induction hypothesis,  $H^2$  has a  $K(\Delta)$ -coloring  $\phi$ . In order to extend  $\phi$  to a  $K(\Delta)$ -coloring of  $G^2$ , it suffices to show that  $w$  has degree at most  $K(\Delta) - 1$  in  $G^2$ . Since  $x$  is adjacent to  $u$ ,  $d_{G^2}(w) \leq d_G(u) + d_G(x) - m(x) - 1$ . From  $m(x) + m(y) \geq d_G(u) - 2$  and  $m(x) \geq m(y)$ , we conclude that  $m(x) \geq \lceil (d_G(u) - 2)/2 \rceil = \lceil d_G(u)/2 \rceil - 1$ . Hence

$$\begin{aligned} d_{G^2}(w) &\leq d_G(u) + d_G(x) - m(x) - 1 \\ &\leq d_G(u) + d_G(x) - \lceil d_G(u)/2 \rceil \\ &\leq \lfloor 3\Delta(G)/2 \rfloor \\ &= K(\Delta) - 1. \end{aligned}$$

**Case 2.**  $x \not\sim u$  and  $y \not\sim u$ .

Similarly as above, let  $H = G - w$ . Then  $H$  is a  $K_4$ -minor free graph with  $\Delta(H) \leq \Delta(G)$ . Note that  $x$  and  $u$  have distance at most 2 in  $H$  since  $d_G(u) \geq 3$  and  $m(x) \geq m(y)$ . Hence  $G^2 - w = H^2$ . It is easy to see that  $m(x) + m(y) = d_G(u)$  and  $d_{G^2}(w) = d_G(u) + d_G(x) - m(x) + 1$ . From  $m(x) \geq m(y)$  we conclude that  $m(x) \geq \lceil d_G(u)/2 \rceil$ . Hence

$$\begin{aligned} d_{G^2}(w) &\leq d_G(u) + d_G(x) - m(x) + 1 \\ &\leq d_G(u) + d_G(x) - \lceil d_G(u)/2 \rceil + 1 \\ &\leq \lfloor 3\Delta(G)/2 \rfloor + 1 \\ &= K(\Delta). \end{aligned}$$

By the induction hypothesis,  $H^2$  has a  $K(\Delta)$ -coloring  $\phi$ , which acts as a partial coloring of  $G^2$  such that all the vertices except  $w$  have been properly colored. Since  $d_{G^2}(u) \leq \Delta(G) + 2 \leq K(\Delta) - 1$ ,  $u$  has at most  $K(\Delta) - 2$  colored neighbors in  $G^2$  in this partial coloring. Hence there are at least two

choices for properly coloring  $u$  before coloring  $w$ . One of these two colorings of  $u$  will imply that the neighbors of  $w$  have at most  $K(\Delta) - 1$  colors. So the coloring  $\phi$  can be extended to a  $K(\Delta)$ -coloring of  $G^2$ .

**Case 3.**  $x \not\sim u$  and  $y \sim u$ .

If  $m(x) = m(y)$ , we may interchange  $x$  and  $y$ , and it falls under Case 1. So assume that  $m(x) > m(y)$ .

First suppose  $d_G(u)$  is odd. Then  $m(x) + m(y) = d_G(u) - 1$  is even.

If  $m(x) \geq m(y) + 4$ , then  $m(x) \geq (d_G(u) + 3)/2$ . This implies that

$$\begin{aligned} d_{G^2}(w) &\leq d_G(u) + d_G(x) - m(x) + 1 \\ &\leq \lfloor 3\Delta(G)/2 \rfloor \\ &= K(\Delta) - 1. \end{aligned}$$

Let  $H = G - w$ . As before,  $H^2 = G^2 - w$  and every  $K(\Delta)$ -coloring of  $H^2$  can be extended to a  $K(\Delta)$ -coloring of  $G^2$ .

If  $m(x) = m(y) + 2$ , we have  $m(x) = (d_G(u) + 1)/2$ . If  $d_G(u) < \Delta(G)$ , then  $d_{G^2}(w) \leq \lfloor 3\Delta(G)/2 \rfloor = K(\Delta) - 1$ , and the same argument applies to this case. Thus we assume that  $d_G(u) = \Delta(G)$ .

If  $x \sim y$ , then  $d_{G^2}(w) \leq d_G(u) + d_G(x) - 1 - (m(x) - 1) \leq \lfloor 3\Delta(G)/2 \rfloor$ . Any  $K(\Delta)$ -coloring of  $G^2 - w$  can be extended to a  $K(\Delta)$ -coloring of  $G^2$ . Assume now that  $x \not\sim y$ . In this case, the calculation shows that  $d_{G^2}(w) \leq K(\Delta)$ . We now show how to choose a  $K(\Delta)$ -coloring of  $G^2 - w$  so that two neighbors of  $w$  in  $G^2$  have the same color. Such a  $K(\Delta)$ -coloring then can be extended to  $w$ .

Since  $d_G(u) = \Delta(G) \geq 5$ , we may choose some  $w' \in M(u, y) \neq \emptyset$ . Let  $H = G - w - w' + xy$ . It is easy to see that  $H$  is also a  $K_4$ -minor free graph with  $\Delta(H) \leq \Delta(G)$ . Moreover,  $G^2 - w - w'$  is a subgraph of  $H^2$ . By the induction hypothesis,  $H^2$  has a  $K(\Delta)$ -coloring  $\phi$ . Note that  $x$  has a color

different from those of  $y$  and all neighbors of  $y$  in  $H$ . Going back to the graph  $G$ , we color  $w'$  with the color  $\phi(x)$ . This coloring can be extended to  $w$  because  $x$  and  $w'$  are neighbors of  $w$  in  $G^2$  with the same color.

Now we consider the case that  $d_G(u)$  is even. Then  $m(x) + m(y) = d_G(u) - 1$  is odd. If  $m(x) \geq m(y) + 3$ , we have  $m(x) \geq d_G(u)/2 + 1$ . If  $d_G(u) < \Delta(G)$ , then

$$\begin{aligned}
d_{G^2}(w) &\leq d_G(u) + d_G(x) - m(x) + 1 \\
&\leq d_G(x) + d_G(u)/2 \\
&\leq \Delta(G) + (\Delta(G) - 1)/2 \\
&\leq \lfloor 3\Delta(G)/2 \rfloor \\
&= K(\Delta) - 1.
\end{aligned}$$

If  $d_G(u) = \Delta(G)$ , the same calculation shows that  $d_{G^2}(w) \leq d_G(x) + d_G(u)/2 \leq \lfloor 3\Delta(G)/2 \rfloor$ . So any  $K(\Delta)$ -coloring of  $G^2 - w$  can be extended to  $G^2$ .

Assume that  $m(x) = m(y) + 1$ . Since  $d_G(u) \geq 4$ ,  $m(y) = d_G(u)/2 - 1 \geq 1$ . Choose  $w' \in M(u, y)$ . Then

$$\begin{aligned}
d_{G^2}(w') &\leq d_G(u) + d_G(y) - m(y) - 1 \\
&\leq d_G(u) + d_G(y) - d_G(u)/2 \\
&\leq \lfloor 3\Delta(G)/2 \rfloor \\
&= K(\Delta) - 1.
\end{aligned}$$

Let  $H = G - w'$ . Then  $H^2 = G^2 - w'$ , and any  $K(\Delta)$ -coloring of  $H^2$  can be extended to  $G^2$ .  $\square$

### 3 Remarks

A graph  $G$  is called *k-degenerate* if every subgraph  $H$  of  $G$  has  $\delta(H) \leq k$ . The *degeneracy* of  $G$  is the minimum  $k$  such that  $G$  is  $k$ -degenerate. The

coloring number  $\text{col}(G)$  of a graph  $G$  is one plus the degeneracy of  $G$ .

In most cases of the proof of Theorem 1, we found a vertex  $w$  of  $G$  such that  $d_{G^2}(w) \leq K(\Delta) - 1$  and  $G^2 - w$  was a subgraph of the square of a  $K_4$ -minor free graph  $H$  with  $|H| = |G| - 1$ . Yet there were two cases where  $d_{G^2}(w) = K(\Delta)$ , in which case  $K(\Delta) = \lfloor 3\Delta(G)/2 \rfloor$ , so the proof actually establishes the following result.

**Theorem 3** *Let  $G$  be a  $K_4$ -minor free graph. Then*

$$\text{col}(G^2) \leq \begin{cases} \Delta(G) + 3 & \text{if } 2 \leq \Delta(G) \leq 3; \\ \lfloor 3\Delta(G)/2 \rfloor + 2 & \text{if } \Delta(G) \geq 4. \end{cases}$$

For  $\Delta(G) = 2$  and  $3$ , the above upper bound for  $\text{col}(G^2)$  is the same as the upper bound for  $\chi(G^2)$  in Theorem 1, hence it is sharp. If  $\Delta(G)$  is an even number  $2k \geq 4$ , then let  $G'_{2k}$  be the graph consisting of vertices  $x_1, x_2, x_3, x_4$  and  $k$  internally disjoint paths of length 2 joining  $x_i$  and  $x_{i+1}$  (addition modulo 4). It is easy to verify that  $G'_{2k}$  is a  $K_4$ -minor free graph,  $\Delta(G'_{2k}) = 2k$ , and  $\delta((G'_{2k})^2) = 3k + 1$ . So the coloring number of  $(G'_{2k})^2$  is at least  $3k + 2$ . If  $\Delta(G)$  is an odd number  $2k + 1 \geq 5$ , then let  $G'_{2k+1}$  be obtained from  $G'_{2k}$  as follows. Replace a length 2 path between  $x_1$  and  $x_2$  by an edge and a length 2 path between  $x_3$  and  $x_4$  by an edge. Then add one length 2 path connecting  $x_1$  and  $x_4$  and one length 2 path connecting  $x_2$  and  $x_3$ . Then  $\Delta(G'_{2k+1}) = 2k + 1$  and  $\delta((G'_{2k+1})^2) = 3k + 2$ . So the coloring number of  $(G'_{2k+1})^2$  is at least  $3k + 3$ . In conclusion, Theorem 3 gives sharp bounds.

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