

# The circular chromatic number of series-parallel graphs with large girth

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## Abstract

It was proved by Hell and Zhu that if  $G$  is a series-parallel graph of girth at least  $2\lfloor(3k-1)/2\rfloor$ , then  $\chi_c(G) \leq 4k/(2k-1)$ . In this paper, we prove that the girth requirement is sharp, i.e., for any  $k \geq 2$ , there is a series-parallel graph  $G$  of girth  $2\lfloor(3k-1)/2\rfloor - 1$  such that  $\chi_c(G) > 4k/(2k-1)$ .

## 1 Introduction

The circular chromatic number (also called the star chromatic number) of a graph is a natural generalization of the notion of chromatic number of a graph. It was introduced by A. Vince [13] in 1988, and has attracted considerable attention since then [1, 2, 6, 7, 8, 11, 14, 18, 19, 20, 21]. Given two integers  $k, d$ , such that  $k \geq d$ , a  $(k, d)$ -coloring of a graph  $G$  is a coloring  $c$  of the vertices of  $G$  with colors  $0, 1, 2, \dots, k-1$  such that for any two adjacent vertices  $x$  and  $y$  of  $G$  we have  $d \leq |c(x) - c(y)| \leq k - d$ . The *circular chromatic number*  $\chi_c(G)$  of  $G$  is defined as the infimum of the ratio  $k/d$  for which there exists a  $(k, d)$ -coloring of  $G$ . Note that any non-trivial graph has circular chromatic number at least 2. It was shown by Vince [13] (cf. also [2, 14] for combinatorial proofs) that if  $G$  is finite then the infimum in this definition is always attained, and hence can be replaced by the minimum.

Note that a  $(k, 1)$ -coloring of a graph  $G$  is just an ordinary  $k$ -coloring of  $G$ . Therefore we have  $\chi_c(G) \leq \chi(G)$  for any graph  $G$ . On the other hand, it was shown in [13] that  $\chi_c(G) > \chi(G) - 1$ . Hence  $\chi(G) = \lceil \chi_c(G) \rceil$ . In this

sense, for a graph  $G$ , the circular chromatic number  $\chi_c(G)$  is a refinement of the chromatic number  $\chi(G)$ , and  $\chi(G)$  is an approximation of  $\chi_c(G)$ .

Since the infimum can be replaced by the minimum, the circular chromatic number of a finite graph is always rational. On the other hand, it was shown in [13] that for any rational number  $r \geq 2$ , there exists a finite graph  $G$  of circular chromatic number  $r$ . Even more is true: it was shown by Zhu [15] that for any integer  $g$  and any rational number  $r \geq 2$ , there exists a graph  $G$  which has girth at least  $g$  and circular chromatic number  $r$  (also see [16] for a discussion of construction of such graphs). This result is a generalization of the result of Erdős concerning the existence of graphs with arbitrarily large girth and chromatic number [5].

A graph  $H$  is called a *minor* of a graph  $G$  if  $H$  is isomorphic to a graph obtained from a subgraph of  $G$  by contracting edges. We say a graph  $G$  is  *$H$ -minor free* if  $H$  is not a minor of  $G$ . It is well-known that if restricted to graphs which are  $H$ -minor free for a fixed graph  $H$ , then the chromatic number will be “small” when the girth is “large”.

If we consider the circular chromatic number instead of the chromatic number, an  $H$ -minor free graph of girth “large enough” will have circular chromatic number “close to” 2: This can be proved by using the following result of Thomassen [12]: *For any finite graph  $H$ , there is an integer  $f(H)$  such that every graph of minimum degree 3 and girth at least  $f(H)$  contains  $H$  as a minor.* Indeed, it can be shown, by induction on the number of vertices, that any  $H$ -minor free graph  $G$  of girth at least  $2kf(H)$  admits a homomorphism to  $C_{2k+1}$ , and hence has circular chromatic number at most  $2 + 1/k$ . Suppose  $G$  is  $H$ -minor free and has girth at least  $2kf(H)$ . Let  $G'$  be obtained from  $G$  by repeatedly deleting all degree 1 vertices (so that  $G'$  has no vertices of degree 1), and let  $G''$  be obtained from  $G'$  by suppressing all vertices of degree 2. Then  $G''$  is  $H$ -minor free and has minimum degree  $\geq 3$ . Hence by the above result of Thomassen,  $G''$  has a cycle of length  $< f(H)$  (unless  $G$  is a tree and hence  $G'$  the empty graph, which is a trivial case). As  $G'$  has girth  $\geq 2kf(H)$ , we conclude that  $G'$  has an induced path  $P$  of length  $2k$ . Now by the induction hypothesis,  $G' - P$  admits a homomorphism to  $C_{2k+1}$ , and such a homomorphism can be easily extended to the induced path  $P$ , and further extended to all the vertices of  $G$ . (In [9], a stronger result was proved: If  $G$  has large girth, then any  $H$ -minor free graph  $G'$  which admits an homomorphism to  $G$  has circular chromatic number close to 2.)

As graphs embeddable in a fixed surface forbids certain finite graph as a minor, it follows that for any surface  $S$ , for any  $\epsilon > 0$ , there is an integer  $g(S, \epsilon)$  such that any graph embeddable in  $S$  and with girth at least  $g(S, \epsilon)$  has circular chromatic number at most  $2 + \epsilon$ .

We note that for any fixed graph  $H$  (respectively, for any fixed surface  $S$ ), and for any  $\epsilon > 0$  the above discussion can be carried out in such a way that actually produces an integer  $g$  such that any  $H$ -minor free graph (respectively, any graph embeddable in  $S$ ) of girth at least  $g$  has circular chromatic number at most  $2 + \epsilon$ . However, such a derived integer  $g$  is usually far from tight. Recently, Hell et al [3] proved that every planar graph of girth at least  $10k - 3$  has circular chromatic number at most  $2 + 1/k$  ( $k \geq 1$ ). This result has been improved by Klostermeyer and Zhang [7] who showed that any planar graph of odd girth at least  $10k - 3$  has circular chromatic number at most  $2 + 1/k$  ( $k \geq 1$ ). However, the result also seems far from sharp, in the sense of girth vs. circular chromatic number of planar graphs.

Given an  $\epsilon > 0$ , it seems to be difficult to find the smallest integer  $g = g(\epsilon)$  such that any  $H$ -minor free graph (or any graph embeddable in a surface  $S$ ) of girth  $\geq g$  has circular chromatic number at most  $2 + \epsilon$ . Nevertheless, we shall prove in this paper that the following result proven in [4] is sharp, in the sense of girth vs. circular chromatic number of  $K_4$ -minor free graphs:

**Theorem 1** ([4]) *If  $G$  is a  $K_4$ -minor free graph of girth at least  $2\lfloor(3k - 1)/2\rfloor$ , then  $\chi_c(G) \leq 4k/(2k - 1)$ .*

To be precise, we shall prove the following result (which maybe viewed as a complement of Theorem 1):

**Theorem 2** *For any  $k \geq 2$ , there is a  $K_4$ -minor free graph  $G$  of girth  $2\lfloor(3k - 1)/2\rfloor - 1$  such that  $\chi_c(G) > 4k/(2k - 1)$ .*

## 2 Preliminaries

It is well known that a graph  $G$  is  $K_4$ -minor free if and only if each block of  $G$  is a series-parallel graph. Therefore to prove Theorem 2, it suffices to deal with series-parallel graphs.

In the remaining of this paper, we assume that  $k \geq 2$  is a fixed integer. Let  $p = \lfloor(3k - 1)/2\rfloor$ . We shall construct a series-parallel graph  $G$  of girth  $2p - 1$  such that  $G$  is not  $(4k, 2k - 1)$ -colorable.

First we recall the following definition from [4]. A two-terminal series-parallel graph  $(G; x, y)$  is defined recursively as follows:

- Let  $V(K_2) = \{0, 1\}$ . Then  $(K_2; 0, 1)$  is a two-terminal series-parallel graph.

- (The parallel construction.) Let  $(G; x, y)$  and  $(G'; x', y')$  be two disjoint two-terminal series-parallel graphs. Define  $G''$  to be the graph obtained from the union of  $G$  and  $G'$  by identifying  $x$  and  $x'$  into a single vertex  $x''$ , and identifying  $y$  and  $y'$  into a single vertex  $y''$ . Then  $(G''; x'', y'')$  is a two-terminal series-parallel graph.
- (The series construction.) Let again  $(G; x, y)$  and  $(G'; x', y')$  be two disjoint two-terminal series-parallel graphs. Define  $G''$  to be the graph obtained from the union of  $G$  and  $G'$  by identifying  $y$  and  $x'$  into a single vertex. Then  $(G''; x, y')$  is a two-terminal series-parallel graph.
- There are no other two-terminal series-parallel graphs.

Note that  $(G; x, y)$  may be a two-terminal series-parallel graph for some pairs  $x, y$  and not for others. A graph  $G$  is a series-parallel graph if there exist some two vertices  $x, y$  such that  $(G; x, y)$  is a two-terminal series-parallel graph. We call the distance in  $G$  between the two terminals  $x, y$ , the *length* of  $(G; x, y)$ . If the two terminals are understood, or are of no significance, we shall write  $G$  for  $(G; x, y)$ .

To prove Theorem 2, we shall construct a series-parallel graph of girth  $2p - 1$  which is not  $(4k, 2k - 1)$ -colorable. To maintain that the constructed series-parallel graphs have girth  $\geq 2p - 1$ , we shall only do the parallel constructions for those pairs of graphs  $G$  and  $G'$  such that the sum of the lengths of the two graphs is at least  $2p - 1$ .

For a two terminal series-parallel graph  $(G; x, y)$  we shall denote by  $L(G)$  the set of colors  $j$  for which there is a  $(4k, 2k - 1)$ -coloring  $c$  of  $G$  such that  $c(x) = 0$  and  $c(y) = j$ . Note that if  $(G; x, y)$  has a  $(4k, 2k - 1)$ -coloring, then there is such a coloring  $c$  with  $c(x) = 0$ . Therefore  $L(G) = \emptyset$  if and only if  $G$  is not  $(4k, 2k - 1)$ -colorable. We shall call  $L(G)$  *the color set* of  $G$ .

Our construction is as follows: Starting from a single edge, we construct a sequence of two-terminal series-parallel graphs  $G$  of girth  $\geq 2p - 1$ . At each step, we shall determine  $L(G)$  for each constructed two-terminal series-parallel graph  $G$ . Once we obtain a two-terminal series-parallel graph  $G$  with  $L(G) = \emptyset$ , then  $G$  is the graph we wanted, i.e.,  $G$  is a series-parallel graph of girth at least  $2p - 1$  (and hence of girth exactly  $2p - 1$ , by Theorem 1) which is not  $(4k, 2k - 1)$ -colorable.

We shall denote by  $I$  the set of all colors, i.e.,  $I = \{0, 1, \dots, 4k - 1\}$ . For two subsets  $A, B$  of  $I$ , we define  $A + B$  as

$$A + B = \{i + j \pmod{4k} : i \in A, j \in B\}.$$

The following three lemmas are straightforward.

**Lemma 1** *If  $G^*$  is obtained from  $G$  and  $G'$  by a series construction, then the length of  $G^*$  is the sum of the lengths of  $G$  and  $G'$ , and that  $L(G^*) = L(G) + L(G')$ .*

**Lemma 2** *If  $G^*$  is obtained from  $G$  and  $G'$  by a parallel construction, then the length of  $G^*$  is equal to the minimum of the lengths of  $G$  and  $G'$ , and that  $L(G^*) = L(G) \cap L(G')$ .*

**Lemma 3**  $L(K_2) = \{2k - 1, 2k, 2k + 1\}$ .

As each two terminal series-parallel graph  $G$  is constructed from  $K_2$  by a sequence of series and parallel constructions, it is straightforward, although tedious sometime, to determine  $L(G)$  for any two-terminal series-parallel graph  $G$ , by using the three lemmas above.

In the remaining, whenever we construct a new series-parallel graph  $G^*$  from two old series-parallel graphs  $G$  and  $G'$  (either by the series construction or the parallel construction), all we care about is the length of  $G^*$  and the color set  $L(G^*)$ . The inner structure of  $G^*$  is of no importance to us. By Lemmas 1 and 2, such information can be obtained from the knowledge of the lengths of  $G$  and  $G'$ , and from  $L(G)$  and  $L(G')$ . Therefore, instead of constructing the series-parallel graph  $G^*$  from  $G$  and  $G'$ , we can simply “construct” the set  $L(G^*)$  from  $L(G)$  and  $L(G')$ , and record the lengths of the corresponding series-parallel graphs.

**Definition 1** *For  $i \geq 1$ , a subset  $S$  of  $I$  is called  $i$ -constructible, if there is a series-parallel graph  $G$  of girth at least  $2p - 1$  which has length  $i$  and for which  $L(G) = S$ .*

By applying Lemmas 1 and 2, we have the following corollary:

**Corollary 1** *Suppose  $A$  is  $i$ -constructible and  $B$  is  $j$ -constructible. Let  $C = A + B$ . Then  $C$  is  $(i + j)$ -constructible. Moreover, if  $i \leq j$  and  $i + j \geq 2p - 1$ , and  $D = A \cap B$ , then  $D$  is  $i$ -constructible.*

As noted above, our goal is to show that the emptyset is  $i$ -constructible for some  $i$  (and hence for all  $i$ ), which means that the corresponding two-terminal series-parallel graph has girth at least  $2p - 1$  and which is not  $(4k, 2k - 1)$ -colorable. We shall repeatedly calculate the sums of two color sets, and the intersections of two color sets. Lemma 4 below, quoted from [4],

illustrates the result of adding two color sets each of which is an “interval” of colors.

We shall view the numbers in  $I$  as cyclically ordered, and all the additions of the elements of  $I$  are carried out modulo  $4k$ , unless otherwise specified. For two elements  $a, b$  of  $I$ , we shall denote by  $[a, b]$  the set  $\{a, a + 1, \dots, b\}$ . For example,  $[2, 5] = \{2, 3, 4, 5\}$  and  $[5, 2] = \{5, 6, \dots, 4k - 1, 0, 1, 2\}$ .

**Lemma 4** *Suppose  $X = [a, b]$  and  $Y = [a', b']$ .*

- *If  $|X| + |Y| \geq 4k + 1$ , then  $X + Y = I$ ;*
- *If  $|X| + |Y| \leq 4k$ , then  $X + Y = [a + a', b + b']$ .*

*Here the sum  $|X| + |Y|$  is the ordinary sum, and the sum  $a + a', b + b'$  are carried out modulo  $4k$ .*

### 3 The proof of Theorem 2

We shall construct, for some integers  $i$ , sequences of smaller and smaller  $i$ -constructible sets, and eventually show that the empty set is  $i$ -constructible for some integer  $i$  (and hence for all  $i$ ). We start with the simplest series-parallel graphs, i.e., paths.

**Lemma 5** *For  $i \geq 1$ , let  $I_i$  be the sets defined as follows:*

1. *If  $i \leq 2k - 1$  is odd, then  $I_i = [2k - i, 2k + i]$ .*
2. *If  $i \leq 2k - 1$  is even, then  $I_i = [4k - i, i]$ .*
3. *For  $i \geq 2k$ ,  $I_i = I$ .*

*Then  $I_i$  is  $i$ -constructible.*

**Proof.** The case  $i = 1$  follows from Lemma 3. For  $i \geq 2$ , it is easy to verify that  $I_i = I_{i-1} + I_1$ . Therefore, by Corollary 1, and by induction on  $i$ ,  $I_i$  is  $i$ -constructible. ■

Indeed,  $I_i$  is just the color set of the path of length  $i$ .

For the remaining part, the construction methods will be slightly different for different values of  $k \pmod{4}$ . We consider four cases.

**Case 1**  $k \equiv 1 \pmod{4}$ .

Assume that  $k = 4n + 1$ .

**Lemma 6** For  $1 \leq m \leq \log_2 n + 1$ , let

1.  $I_{6n+1}^{(m)} = [2n + 2^{m-1}, (6n + 2)] \cup [(10n + 2), (14n + 4) - 2^{m-1}]$ ,
2.  $I_{6n}^{(m)} = [2n + 2^{m-1}, (6n + 1) - 2^{m-1}] \cup [(10n + 3) + 2^{m-1}, (14n + 4) - 2^{m-1}]$ ,
3.  $I_{4n}^{(m)} = [2^m - 1, (4n + 1) - 2^m] \cup [(12n + 3) + 2^m, (16n + 5) - 2^m]$ ,
4.  $I_{4n+1}^{(m)} = [4n + 2^m, (8n + 2) - 2^m] \cup [(8n + 2) + 2^m, (12n + 4) - 2^m]$ ,
5.  $I_{8n+1}^{(m)} = [2^{m+1} - 1, (16n + 5) - 2^{m+1}]$ ,
6.  $I_{8n}^{(m)} = [2^{m+1} - 1, (8n + 2) - 2^{m+1}] \cup [(8n + 2) + 2^{m+1}, (16n + 5) - 2^{m+1}]$ .

The sets  $I_i^{(j)}$  defined above are  $i$ -constructible.

**Proof.** We prove Lemma 6 by induction on  $m$ . Note that since  $k = 4n + 1$ , we have  $p = 6n + 1$ . We can do parallel construction to two two-terminal series-parallel graphs  $G$  and  $G'$  if and only if the sum of the lengths of  $G$  and  $G'$  is at least  $2p - 1 = 12n + 1$ .

If  $m = 1$ , let  $A = I_{4n} \cap I_{8n+1}$ , then straightforward calculation shows that

$$A = [1, 4n] \cup [12n + 4, 16n + 3].$$

Since  $I_{4n}$  and  $I_{8n+1}$  are  $4n$ -constructible and  $(8n + 1)$ -constructible, respectively (by Lemma 5), and since  $4n + 8n + 1 = 12n + 1 = 2p - 1$ , it follows from Corollary 1 that  $A$  is  $4n$ -constructible.

Let  $B = I_{4n+1} \cap I_{8n}$ , then

$$B = [4n + 1, 8n] \cup [8n + 4, 12n + 3],$$

and similarly as above,  $B$  is  $(4n + 1)$ -constructible.

Let  $I_{6n+1}^{(1)} = I_{6n+1} \cap I_{6n+2}$ , then straightforward calculation shows that

$$I_{6n+1}^{(1)} = [2n + 1, 6n + 2] \cup [10n + 2, 14n + 3],$$

and similarly as above,  $I_{6n+1}^{(1)}$  is  $(6n + 1)$ -constructible.

Let  $I_{6n}^{(1)} = I_{6n} \cap I_{6n+1}$ , then

$$I_{6n}^{(1)} = [2n + 1, 6n] \cup [10n + 4, 14n + 3],$$

and similarly,  $I_{6n}^{(1)}$  is  $6n$ -constructible.

Let  $C = I_{6n}^{(1)} + I_{6n}^{(1)}$ , then

$$C = [4n + 2, 12n + 2] \cup [12n + 5, 4n - 1],$$

and  $C$  is  $12n$ -constructible.

Let  $I_{4n}^{(1)} = A \cap C$ , then

$$I_{4n}^{(1)} = [1, 4n - 1] \cup [12n + 5, 16n + 3],$$

and  $I_{4n}^{(1)}$  is  $4n$ -constructible (because  $A$  is  $4n$ -constructible,  $C$  is  $12n$ -constructible, and  $4n + 12n \geq 2p - 1$ ).

Let  $I_{4n+1}^{(1)} = B \cap C$ , then

$$I_{4n+1}^{(1)} = [4n + 2, 8n] \cup [8n + 4, 12n + 2],$$

and similarly as above,  $I_{4n+1}^{(1)}$  is  $(4n + 1)$ -constructible.

Let  $I_{8n+1}^{(1)} = (I_{4n}^{(1)} + I_{4n+1}^{(1)})$ , then

$$I_{8n+1}^{(1)} = [3, 16n + 1],$$

and  $I_{8n+1}^{(1)}$  is  $(8n+1)$ -constructible, because  $I_{4n}^{(1)} + I_{4n+1}^{(1)}$  is  $(8n+1)$ -constructible (by Corollary 1) and  $I_{8n+1}$  is also  $(8n + 1)$ -constructible.

Let  $I_{8n}^{(1)} = (I_{4n}^{(1)} + I_{4n}^{(1)}) \cap I_{8n+1}^{(1)}$ , then

$$I_{8n}^{(1)} = [3, 8n - 2] \cup [8n + 6, 16n + 1],$$

and similarly as above,  $I_{8n}^{(1)}$  is  $8n$ -constructible.

This completes the proof of the  $m = 1$  case of Lemma 6.



Suppose Lemma 6 is true for  $m = i$ , i.e., the sets  $I_{6n+1}^{(i)}, I_{6n}^{(i)}, I_{4n}^{(i)}, I_{4n+1}^{(i)}, I_{8n+1}^{(i)}, I_{8n}^{(i)}$  defined as in Lemma 6 are  $(6n+1)$ -,  $6n$ -,  $4n$ -,  $(4n+1)$ -,  $(8n+1)$ - $8n$ -constructible, respectively. Assume  $i+1 \leq \log_2 n + 1$ . We shall use the induction hypothesis to prove that  $I_{6n+1}^{(i+1)}, I_{6n}^{(i+1)}, I_{4n}^{(i+1)}, I_{4n+1}^{(i+1)}, I_{8n+1}^{(i+1)}, I_{8n}^{(i+1)}$  defined as in Lemma 6 are  $(6n+1)$ -,  $6n$ -,  $12n$ -,  $4n$ -,  $(4n+1)$ -,  $(8n+1)$ -,  $8n$ -constructible, respectively.

By the induction hypothesis,

$$I_{4n}^{(i)} = [2^i - 1, (4n+1) - 2^i] \cup [(12n+3) + 2^i, (16n+5) - 2^i]$$

$$I_{8n+1}^{(i)} = [2^{i+1} - 1, (16n+5) - 2^{i+1}]$$

are  $4n$ - and  $(8n+1)$ -constructible, respectively. Let  $X = I_{4n}^{(i)} \cap I_{8n+1}^{(i)}$ . Since  $i+1 \leq \log_2 n + 1$ , we have  $2^i \leq n$ . This implies that

$$((4n+1) - 2^i) - (2^{i+1} - 1) = (4n+2) - 3 \times 2^i \geq 0,$$

$$((16n+5) - 2^{i+1}) - ((12n+3) + 2^i) = (4n+2) - 3 \times 2^i \geq 0.$$

So

$$2^i - 1 \leq 2^{i+1} - 1 \leq (4n+1) - 2^i,$$

$$(12n+3) + 2^i \leq (16n+5) - 2^{i+1} \leq (16n+5) - 2^i.$$

Therefore

$$X = I_{4n}^{(i)} \cap I_{8n+1}^{(i)} = [2^{i+1} - 1, (4n+1) - 2^i] \cup [(12n+3) + 2^i, (16n+5) - 2^{i+1}].$$

Similarly,

$$Y = I_{4n+1}^{(i)} \cap I_{8n}^{(i)} = [4n+2^i, (8n+2) - 2^{i+1}] \cup [(8n+2) + 2^{i+1}, (12n+4) - 2^i].$$

By Corollary 1,  $X$  and  $Y$  are  $4n$ -,  $(4n+1)$ -constructible, respectively.

Let  $I_{6n+1}^{(i+1)} = (Y + I_{2n}) \cap I_{6n+1}^{(i)}$ . Firstly, by applying lemma 4, we have

$$Y + I_{2n} = [2n + 2^i, (10n+2) - 2^{i+1}] \cup [(6n+2) + 2^{i+1}, (14n+4) - 2^i].$$

Since  $((10n+2) - 2^{i+1}) - ((6n+2) + 2^{i+1}) = 4n - 2^{i+2} \geq 0$  (as  $i+1 \leq \log_2 n + 1$ ), the two intervals in the formula above intersect, and hence their union is a single interval, i.e.,

$$Y + I_{2n} = [2n + 2^i, (14n+4) - 2^i].$$

Similarly as the argument in the previous paragraph, we conclude that

$$(Y + I_{2n}) \cap I_{6n+1}^{(i)} = [2n + 2^i, (6n+2)] \cup [(10n+2), (14n+4) - 2^i].$$

By Corollary 1,  $I_{6n+1}^{(i+1)} = (Y + I_{2n}) \cap I_{6n+1}^{(i)}$  is  $(6n + 1)$ -constructible.

Let  $I_{6n}^{(i+1)} = (X + I_{2n}) \cap I_{6n+1}^{(i+1)}$ . Then by the same argument as above

$$I_{6n}^{(i+1)} = [2n + 2^i, (6n + 1) - 2^i] \cup [(10n + 3) + 2^i, (14n + 4) - 2^i]$$

and  $I_{6n}^{(i+1)}$  is  $6n$ -constructible, because  $X + I_{2n}$  is  $6n$ -constructible, and  $I_{6n+1}^{(i+1)}$  is  $(6n + 1)$ -constructible.

Let  $Z = I_{6n}^{(i+1)} + I_{6n}^{(i+1)}$ . Then

$$Z = [4n + 2^{i+1}, (12n + 4) - 2^{i+1}] \cup [(12n + 3) + 2^{i+1}, (4n + 1) - 2^{i+1}]$$

and  $Z$  is  $12n$ -constructible by Corollary 1.

Let  $I_{4n}^{(i+1)} = X \cap Z$ . Then

$$I_{4n}^{(i+1)} = [2^{i+1} - 1, (4n + 1) - 2^{i+1}] \cup [(12n + 3) + 2^{i+1}, (16n + 5) - 2^{i+1}]$$

and  $I_{4n}^{(i+1)}$  is  $4n$ -constructible, because  $X$  is  $4n$ -constructible, and  $Z$  is  $12n$ -constructible.

Let  $I_{(4n+1)}^{(i+1)} = Y \cap Z$ . Then

$$I_{(4n+1)}^{(i+1)} = [4n + 2^{i+1}, (8n + 2) - 2^{i+1}] \cup [(8n + 2) + 2^{i+1}, (12n + 4) - 2^{i+1}]$$

and  $I_{(4n+1)}^{(i+1)}$  is  $(4n + 1)$ -constructible, because  $Y$  is  $(4n + 1)$ -constructible, and  $Z$  is  $12n$ -constructible.

Let  $I_{8n+1}^{(i+1)} = I_{4n+1}^{(i+1)} + I_{4n}^{(i+1)}$ . Then

$$I_{8n+1}^{(i+1)} = [2^{i+2} - 1, (16n + 5) - 2^{i+2}]$$

and  $I_{8n+1}^{(i+1)}$  is  $(8n + 1)$ -constructible, because  $I_{4n+1}^{(i+1)} + I_{4n}^{(i+1)}$  and  $I_{8n+1}^{(i)}$  are both  $(8n + 1)$ -constructible.

Let  $I_{8n}^{(i+1)} = (I_{4n}^{(i+1)} + I_{4n}^{(i+1)}) \cap I_{8n+1}^{(i+1)}$ . Then

$$I_{8n}^{(i+1)} = [2^{i+2} - 1, (8n + 2) - 2^{i+2}] \cup [(8n + 2) + 2^{i+2}, (16n + 5) - 2^{i+2}]$$

and  $I_{8n}^{(i+1)}$  is  $8n$ -constructible, because  $I_{4n}^{(i+1)} + I_{4n}^{(i+1)}$  and  $I_{8n}^{(i)}$  are both  $8n$ -constructible, and  $I_{8n+1}^{(i+1)}$  is  $(8n + 1)$ -constructible. This completes the proof of Lemma 6. ■

**Case 2**  $k \equiv 2 \pmod{4}$ .

Assume that  $k = 4n + 2$ .

**Lemma 7** For  $1 \leq m \leq \log_2 n + 1$ , let

1.  $I_{6n+2}^{(m)} = [(2n + 1), (6n + 3) - 2^{m-1}] \cup [(10n + 5) + 2^{m-1}, (14n + 7)]$ ,
2.  $I_{6n+1}^{(m)} = [(2n + 2) + 2^{m-1}, (6n + 3) - 2^{m-1}] \cup [(10n + 5) + 2^{m-1}, (14n + 6) - 2^{m-1}]$ ,
3.  $I_{4n}^{(m)} = [2^{m-1}, (4n + 1) - 2^m] \cup [(12n + 7) + 2^m, (16n + 8) - 2^{m-1}]$ ,
4.  $I_{4n+1}^{(m)} = [(4n + 2) + 2^m, (8n + 4) - 2^m] \cup [(8n + 4) + 2^m, (12n + 6) - 2^m]$ ,
5.  $I_{8n+3}^{(m)} = [2^m, (16n + 8) - 2^m]$ ,
6.  $I_{8n+2}^{(m)} = [2^m, (8n + 4) - 2^{m+1}] \cup [(8n + 4) + 2^{m+1}, (16n + 8) - 2^m]$ .

The sets  $I_i^{(j)}$  defined above are *i-constructible*.

**Proof.** The proof of Lemma 7 is similar to the proof of Lemma 6. We shall only give the rules of how to construct the sets  $I_i^{(1)}$ , and how to construct the sets  $I_i^{(j+1)}$  from the sets  $I_i^{(j)}$ . Since  $k = 4n + 2$ , we have  $p = 6n + 2$ , hence  $2p - 1 = 12n + 3$ . Therefore, we can do parallel construction to two-terminal series-parallel graphs  $G$  and  $G'$  if and only if the sum of the lengths of the two graphs is at least  $12n + 3$ .

The sets  $I_i^{(1)}$  are constructed as follows:

1.  $I_{6n+2}^{(1)} = I_{6n+2} \cap I_{6n+3}$ ,
2.  $I_{6n+1}^{(1)} = I_{6n+1} \cap I_{6n+2}$ ,
3.  $I_{4n}^{(1)} = (I_{6n+1}^{(1)} + I_{6n+1}^{(1)}) \cap I_{4n} \cap I_{8n+3}$ ,
4.  $I_{4n+1}^{(1)} = (I_{6n+1}^{(1)} + I_{6n+1}^{(1)}) \cap I_{4n+1} \cap I_{8n+2}$ ,

5.  $I_{8n+3}^{(1)} = I_{4n+1}^{(1)} + I_{4n+2}$ ,
6.  $I_{8n+2}^{(1)} = (I_{4n+1}^{(1)} + I_{4n+1}^{(1)}) \cap I_{8n+3}^{(1)}$ .

The sets  $I_i^{(j+1)}$  are constructed from the sets  $I_i^{(j)}$  as follows:

1.  $I_{6n+2}^{(j+1)} = ((I_{4n+1}^{(j)} \cap I_{8n+2}^{(j)}) + I_{2n+1}) \cap I_{6n+2}^{(j)}$ ,
2.  $I_{6n+1}^{(j+1)} = ((I_{4n}^{(j)} \cap I_{8n+3}^{(j)}) + I_{2n+1}) \cap I_{6n+2}^{(j+1)}$ ,
3.  $I_{4n}^{(j+1)} = (I_{6n+1}^{(j+1)} + I_{6n+1}^{(j+1)}) \cap I_{4n}^{(j)} \cap I_{8n+3}^{(j)}$ ,
4.  $I_{4n+1}^{(j+1)} = (I_{6n+1}^{(j+1)} + I_{6n+1}^{(j+1)}) \cap I_{4n+1}^{(j)} \cap I_{8n+2}^{(j)}$ ,
5.  $I_{8n+3}^{(j+1)} = I_{4n+1}^{(j+1)} + I_{4n+2}$ ,
6.  $I_{8n+2}^{(j+1)} = (I_{4n+1}^{(j+1)} + I_{4n+1}^{(j+1)}) \cap I_{8n+3}^{(j+1)}$ .

■

**Case 3**  $k \equiv 3 \pmod{4}$ .

Assume that  $k = 4n + 3$ .

**Lemma 8** For  $1 \leq m \leq \log_2(n+1) + 1$ , let

1.  $I_{6n+4}^{(m)} = [(2n+1), (6n+5) - 2^{m-1}] \cup [(10n+7) + 2^{m-1}, (14n+11)]$ ,
2.  $I_{6n+3}^{(m)} = [(2n+2) + 2^{m-1}, (6n+5) - 2^{m-1}] \cup [(10n+7) + 2^{m-1}, (14n+10) - 2^{m-1}]$ ,
3.  $I_{4n+2}^{(m)} = [2^m - 1, (4n+3) - 2^m] \cup [(12n+9) + 2^m, (16n+13) - 2^m]$ ,

4.  $I_{4n+3}^{(m)} = [(4n+2) + 2^m, (8n+6) - 2^m] \cup [(8n+6) + 2^m, (12n+10) - 2^m]$ ,
5.  $I_{8n+5}^{(m)} = [2^{m+1} - 1, (16n+13) - 2^{m+1}]$ ,
6.  $I_{8n+4}^{(m)} = [2^{m+1} - 1, (8n+6) - 2^{m+1}] \cup [(8n+6) + 2^{m+1}, (16n+13) - 2^{m+1}]$ .

The sets  $I_i^{(j)}$  defined above are  $i$ -constructible.

**Proof.** Similarly, we shall only give the rules of how to construct the sets  $I_i^{(1)}$ , and how to construct the sets  $I_i^{(j+1)}$  from the sets  $I_i^{(j)}$ . Since  $k = 4n + 3$ , we have  $p = 6n + 4$ , hence  $2p - 1 = 12n + 7$ . Therefore, we can do parallel construction to two-terminal series-parallel graphs  $G$  and  $G'$  if and only if the sum of the lengths of the two graphs is at least  $12n + 7$ .

The sets  $I_i^{(1)}$  are constructed as follows:

1.  $I_{6n+4}^{(1)} = I_{6n+4} \cap I_{6n+5}$ ,
2.  $I_{6n+3}^{(1)} = I_{6n+3} \cap I_{6n+4}$ ,
3.  $I_{4n+2}^{(1)} = (I_{6n+3}^{(1)} + I_{6n+3}^{(1)}) \cap I_{4n+2} \cap I_{8n+5}$ ,
4.  $I_{4n+3}^{(1)} = (I_{6n+3}^{(1)} + I_{6n+3}^{(1)}) \cap I_{4n+3} \cap I_{8n+4}$ ,
5.  $I_{8n+5}^{(1)} = I_{4n+2}^{(1)} + I_{4n+3}^{(1)}$ ,
6.  $I_{8n+4}^{(1)} = (I_{4n+2}^{(1)} + I_{4n+2}^{(1)}) \cap I_{8n+5}^{(1)}$ .

The sets  $I_i^{(j+1)}$  are constructed from the sets  $I_i^{(j)}$  as follows:

1.  $I_{6n+4}^{(j+1)} = ((I_{4n+3}^{(j)} \cap I_{8n+4}^{(j)}) + I_{2n+1}) \cap I_{6n+4}^{(j)}$ ,
2.  $I_{6n+3}^{(j+1)} = ((I_{4n+2}^{(j)} \cap I_{8n+5}^{(j)}) + I_{2n+1}) \cap I_{6n+4}^{(j+1)}$ ,

3.  $I_{4n+2}^{(j+1)} = (I_{6n+3}^{(j+1)} + I_{6n+3}^{(j+1)}) \cap I_{4n+2}^{(j)} \cap I_{8n+5}^{(j)}$ ,
4.  $I_{4n+3}^{(j+1)} = (I_{6n+3}^{(j+1)} + I_{6n+3}^{(j+1)}) \cap I_{4n+3}^{(j)} \cap I_{8n+4}^{(j)}$ ,
5.  $I_{8n+5}^{(j+1)} = I_{4n+2}^{(j+1)} + I_{4n+3}^{(j+1)}$ ,
6.  $I_{8n+4}^{(j+1)} = (I_{4n+2}^{(j+1)} + I_{4n+2}^{(j+1)}) \cap I_{8n+5}^{(j+1)}$ .

■

**Case 4**  $k \equiv 0 \pmod{4}$ .

Assume that  $k = 4n$ .

**Lemma 9** For  $1 \leq m \leq \log_2 n + 1$ , let

1.  $I_{6n-1}^{(m)} = [2n + 2^{m-1}, 6n] \cup [10n, 14n - 2^{m-1}]$ ,
2.  $I_{6n-2}^{(m)} = [2n + 2^{m-1}, (6n - 1) - 2^{m-1}] \cup [(10n + 1) + 2^{m-1}, 14n - 2^{m-1}]$ ,
3.  $I_{4n-2}^{(m)} = [2^{m-1}, (4n - 1) - 2^m] \cup [(12n + 1) + 2^m, 16n - 2^{m-1}]$ ,
4.  $I_{4n-1}^{(m)} = [4n + 2^m, 8n - 2^m] \cup [8n + 2^m, 12n - 2^m]$ ,
5.  $I_{8n-1}^{(m)} = [2^m, 16n - 2^m]$ ,
6.  $I_{8n-2}^{(m)} = [2^m, 8n - 2^{m+1}] \cup [8n + 2^{m+1}, 16n - 2^m]$ .

The sets  $I_i^{(j)}$  defined above are  $i$ -constructible.

**Proof.** Similarly, we shall only give the rules of how to construct the sets  $I_i^{(1)}$ , and how to construct the sets  $I_i^{(j+1)}$  from the sets  $I_i^{(j)}$ . Since  $k = 4n$ , we have  $p = 6n - 1$ , hence  $2p - 1 = 12n - 3$ . Therefore, we can do parallel construction to two-terminal series-parallel graphs  $G$  and  $G'$  if and only if the sum of the lengths of the two graphs is at least  $12n - 3$ .

The sets  $I_i^{(1)}$  are constructed as follows:

1.  $I_{6n-1}^{(1)} = I_{6n-1} \cap I_{6n},$
2.  $I_{6n-2}^{(1)} = I_{6n-2} \cap I_{6n-1},$
3.  $I_{4n-2}^{(1)} = (I_{6n-2}^{(1)} + I_{6n-2}^{(1)}) \cap I_{4n-2} \cap I_{8n-1},$
4.  $I_{4n-1}^{(1)} = (I_{6n-2}^{(1)} + I_{6n-2}^{(1)}) \cap I_{4n-1} \cap I_{8n-2},$
5.  $I_{8n-1}^{(1)} = I_{4n-1}^{(1)} + I_{4n},$
6.  $I_{8n-2}^{(1)} = (I_{4n-1}^{(1)} + I_{4n-1}^{(1)}) \cap I_{8n-1}^{(1)}.$

The sets  $I_i^{(j+1)}$  are constructed from the sets  $I_i^{(j)}$  as follows:

1.  $I_{6n-1}^{(j+1)} = ((I_{4n-1}^{(j)} \cap I_{8n-2}^{(j)}) + I_{2n}) \cap I_{6n-1}^{(j)},$
2.  $I_{6n-2}^{(j+1)} = ((I_{4n-2}^{(j)} \cap I_{8n-1}^{(j)}) + I_{2n}) \cap I_{6n-1}^{(j+1)},$
3.  $I_{4n-2}^{(j+1)} = (I_{6n-2}^{(j+1)} + I_{6n-2}^{(j+1)}) \cap I_{4n-2}^{(j)} \cap I_{8n-1}^{(j)},$
4.  $I_{4n-1}^{(j+1)} = (I_{6n-2}^{(j+1)} + I_{6n-2}^{(j+1)}) \cap I_{4n-1}^{(j)} \cap I_{8n-2}^{(j)},$
5.  $I_{8n-1}^{(j+1)} = I_{4n-1}^{(j+1)} + I_{4n},$
6.  $I_{8n-2}^{(j+1)} = (I_{4n-1}^{(j+1)} + I_{4n-1}^{(j+1)}) \cap I_{8n-1}^{(j+1)}.$

■

**Proof of Theorem 2** In each of the four cases above, as  $m$  becomes larger and larger, the  $i$ -constructible sets become smaller and smaller. This process can continue until for some  $i$ , the emptyset becomes  $i$ -constructible. In the statement of the four lemmas above, an upper bound for  $m$  is given. This is merely because of the difficulty in expressing the emptyset as an interval of

the form  $[a, b]$ . Indeed, by our convention, for any  $a, b$ , the interval  $[a, b]$  is always nonempty.

Consider the case  $k = 4n + 1$ . Suppose we have carried out the inductive constructions of the sets to the  $m$ th step, where  $m = \lfloor \log_2 n + 1 \rfloor$  (which implies that  $n + 1 \leq 2^m \leq 2n$ , and hence  $2n + 2 \leq 2^{m+1} \leq 4n$ ). Now we may continue in the same way, and try to carry out one more induction step. However, in doing so we shall indeed conclude that the empty set is  $i$ -constructible for some  $i$ .

Let us follow the construction described in the proof of Lemma 6 for one more step. Let  $X = I_{4n}^{(m)} \cap I_{8n+1}^{(m)}$ . Then

$$X = [2^{m+1} - 1, (4n + 1) - 2^m] \cup [(12n + 3) + 2^m, (16n + 5) - 2^{m+1}]$$

is  $4n$ -constructible. Let  $Y = I_{4n+1}^{(m)} \cap I_{8n}^{(m)}$ ,  $I_{6n+1}^{(m+1)} = (Y + I_{2n}) \cap I_{6n}^{(m)}$ ,  $I_{6n}^{(m+1)} = (X + I_{2n}) \cap I_{6n+1}^{(m+1)}$ , and  $Z = I_{6n}^{(m+1)} + I_{6n}^{(m+1)}$ . Then  $Y$ ,  $I_{6n+1}^{(m+1)}$ ,  $I_{6n}^{(m+1)}$  and  $Z$  are  $(4n + 1)$ -,  $(6n + 1)$ -,  $6n$ - and  $12n$ -constructible, respectively. The same calculation shows that

$$Z = [4n + 2^{m+1}, (12n + 4) - 2^{m+1}] \cup [(12n + 3) + 2^{m+1}, (4n + 1) - 2^{m+1}].$$

Let  $I_{4n}^{(m+1)} = X \cap Z$ . Then  $I_{4n}^{(m+1)}$  is  $4n$ -constructible. However, by noting that  $2^{m+1} \geq 2n + 2$ , we can conclude that  $X \cap Z = \emptyset$ . Therefore the empty set is  $4n$ -constructible. As observed earlier, this means that the corresponding series-parallel graph is not  $(4k, 2k - 1)$ -colorable.

The other cases are similar, and we omit the details. ■

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